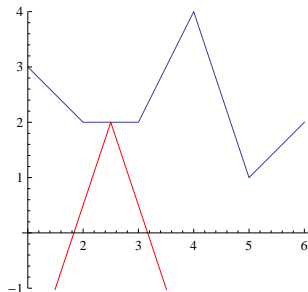


Pijavskij algoritam С. А. Пиявский, Один алгоритм отыскания абсолютного экстремума функции, Ж. вычисл. матем. и матем. физ. 12(1972), 888-896

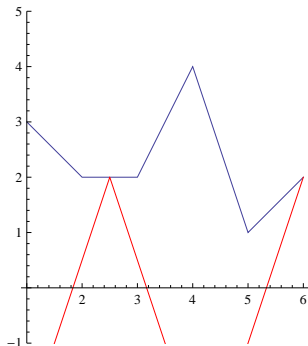
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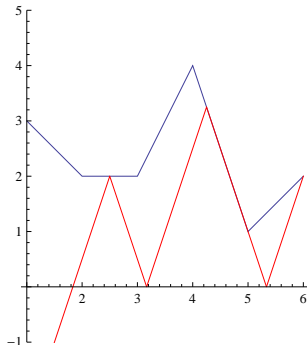
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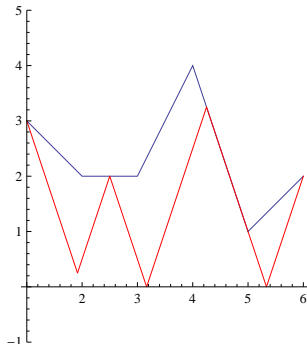
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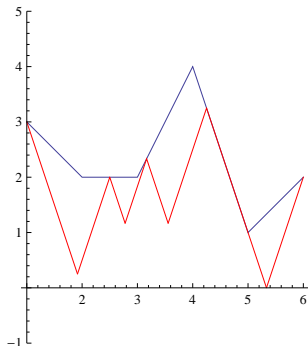
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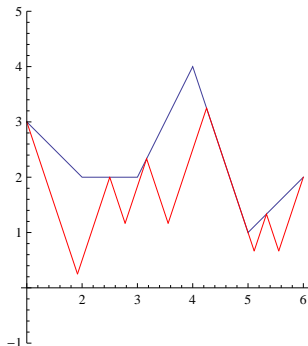
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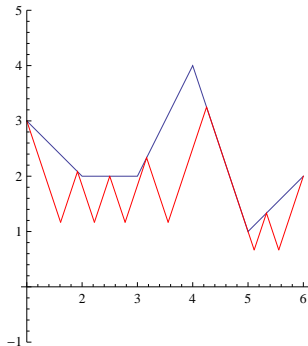
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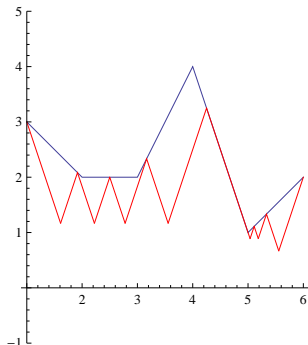
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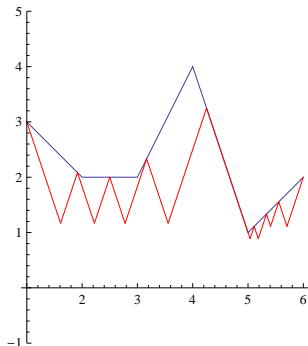
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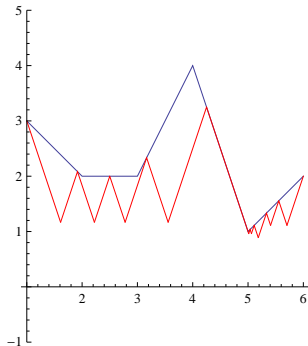
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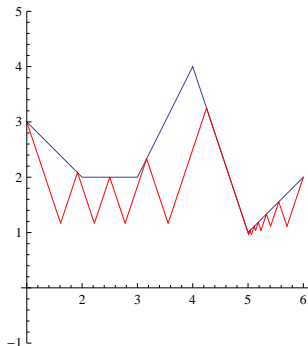
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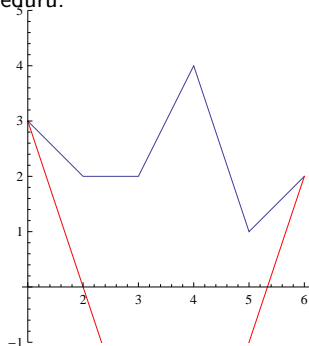
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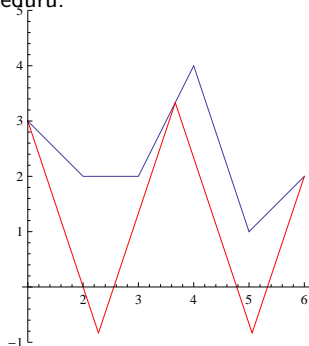
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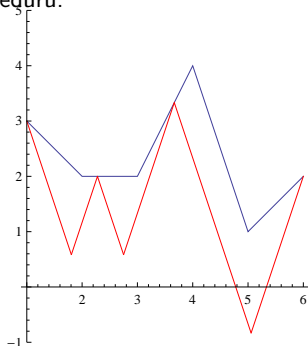
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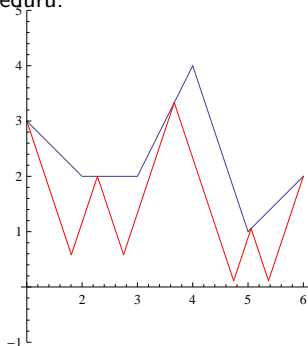
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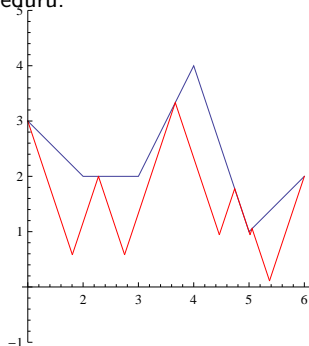
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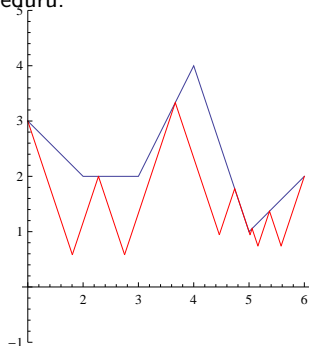
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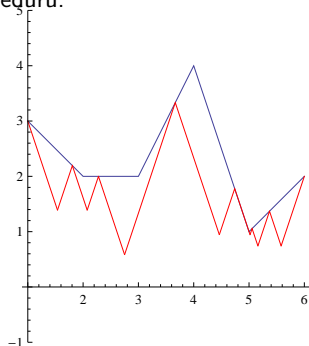
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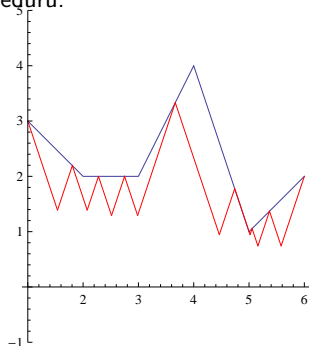
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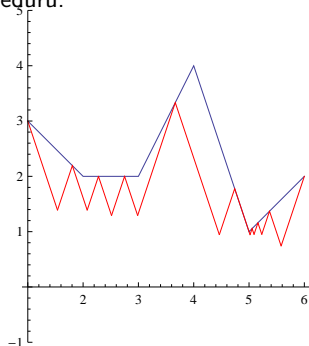
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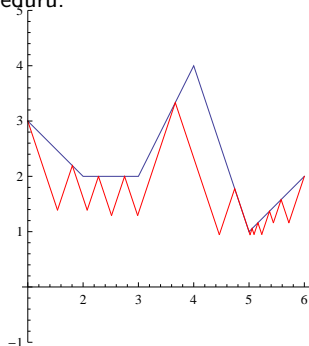
Shubert algoritam (B. Shubert, *A sequential method seeking the global maximum of a function*, SIAM Journal on Numerical Analysis, 9(1972), 379–388)

- Odrediti pravce $x \mapsto f(a) - L(x - a)$ i $x \mapsto f(b) + L(x - b)$ i točku njihovog sjecišta ($U(a, b, f, L)$, $B(a, b, f, L)$)
- Staviti $u_1 = U(a, b, f, L)$. Od dva intervala $[a, u_1]$, $[u_1, b]$ izabrati onaj s manjom B -vrijednosti.
- Ako je $[a, u_1]$ izabrani interval, staviti $u_2 = U(a, u_1, f, L)$.
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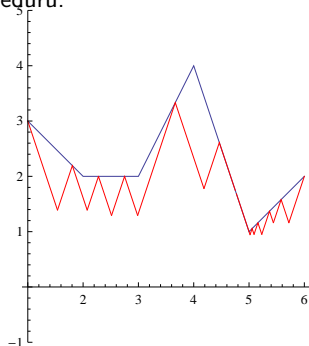
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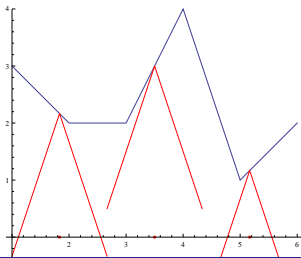
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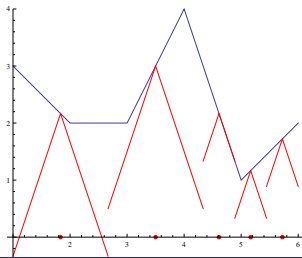
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Izračunati \mathcal{B} -vrijednosti za tri nova podintervala.
Staviti $n := n + 2$ i prijeći na **K1**;



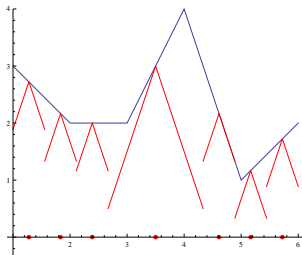
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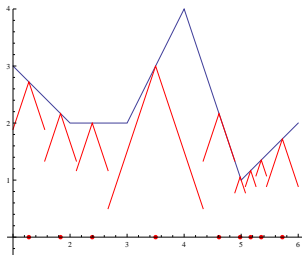
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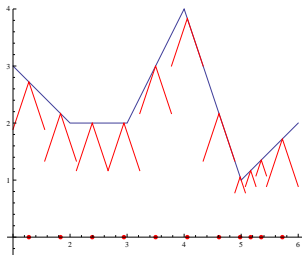
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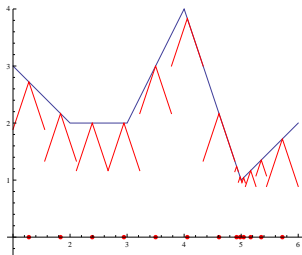
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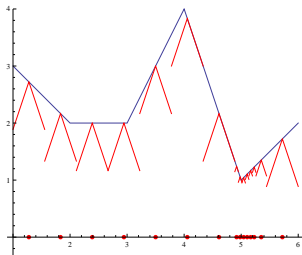
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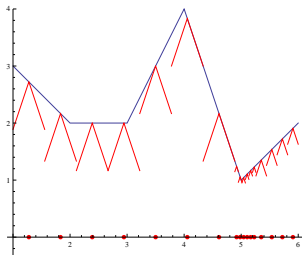
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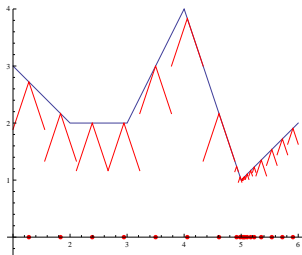
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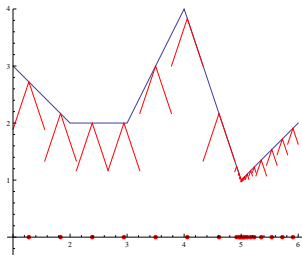
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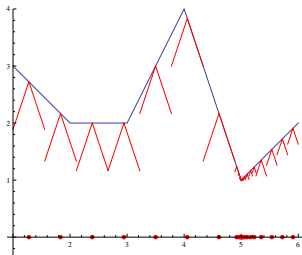
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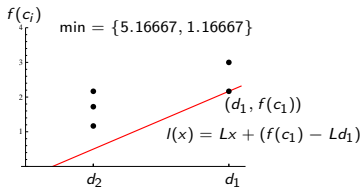
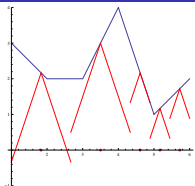
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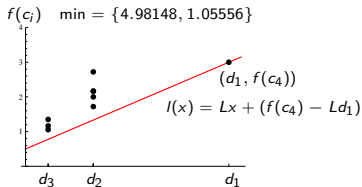
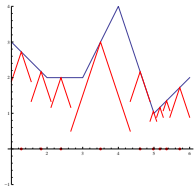


Traženje trenutno optimalnog intervala ($d_1 = \frac{b-a}{6}$, $d_2 = \frac{b-a}{18}$, $d_3 = \frac{b-a}{54}$, $d_4 = \frac{b-a}{162}$)

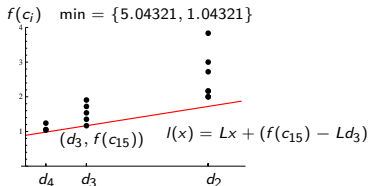
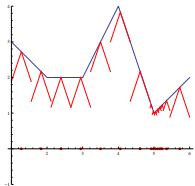
2. iteracija



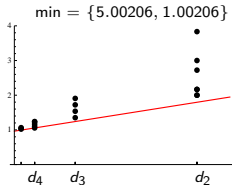
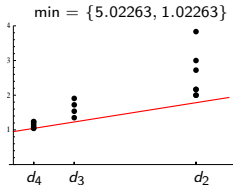
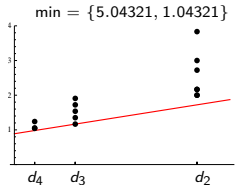
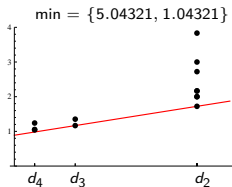
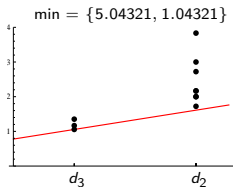
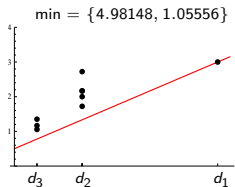
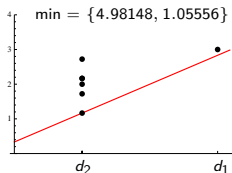
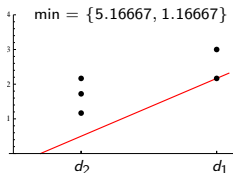
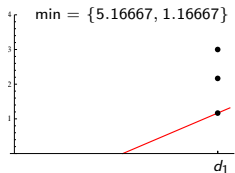
4. iteracija



7. iteracija

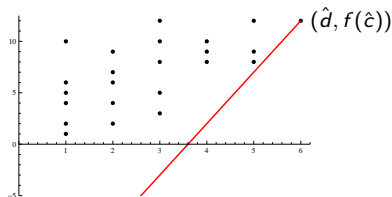


Traženje trenutno optimalnog intervala ($d_1 = \frac{b-a}{6}$, $d_2 = \frac{b-a}{18}$, $d_3 = \frac{b-a}{54}$, $d_4 = \frac{b-a}{162}$)



Za $\hat{L} = 5$ sve točke leže iznad pravca

$$l_1(x) = \hat{L}x + (f(\hat{c}) - \hat{L}\hat{d}) = 5x + (12 - 5 \cdot 6)$$



$$f(\hat{c}) - \hat{L}\hat{d} \leq f(c_i) - \hat{L}d_i \quad \text{za sve } i$$

Dakle, točka $\hat{T} = (\hat{d}, f(\hat{c}))$ reprezentira interval na kome se postiže najmanja \mathcal{B} -vrijednost.

Traženje trenutno optimalnog intervala ekvivalentno je traženju točke $T_i = (d_i, f(c_i))$, $i = 1, \dots, n$ kojom prolazi pravac s koeficijentom smjera $L > 0$ i najmanjim odsječkom na ordinati.

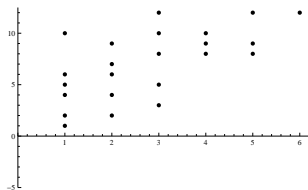
Traženje potencijalno optimalnih intervala

Međutim, kao što se može vidjeti ima i boljih pravaca s većim koeficijentom smjera, a koji imaju spomenuto svojstvo:

$$f(c_1) - L_1 d_1 \leq f(c_i) - L_1 d_i \quad \text{za sve } i$$

$$f(c_2) - L_2 d_2 \leq f(c_i) - L_2 d_i \quad \text{za sve } i$$

$$f(c_3) - L_2 d_3 \leq f(c_i) - L_2 d_i \quad \text{za sve } i$$



- svaki od navedenih pravaca određuje barem jednu potencijalno optimalnu točku;
- sve dobre točke leže na konveksnoj ljusci;

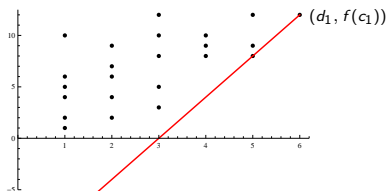
Traženje potencijalno optimalnih intervala

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$$f(c_1) - L_1 d_1 \leq f(c_i) - L_1 d_i \quad \text{za sve } i$$

$$f(c_2) - L_2 d_2 \leq f(c_i) - L_2 d_i \quad \text{za sve } i$$

$$f(c_3) - L_2 d_3 \leq f(c_i) - L_2 d_i \quad \text{za sve } i$$



- svaki od navedenih pravaca određuje barem jednu potencijalno optimalnu točku;
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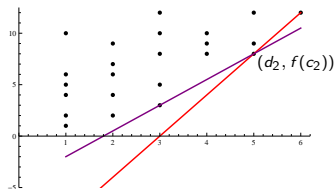
Traženje potencijalno optimalnih intervala

Međutim, kao što se može vidjeti ima i boljih pravaca s većim koeficijentom smjera, a koji imaju spomenuto svojstvo:

$$f(c_1) - L_1 d_1 \leq f(c_i) - L_1 d_i \quad \text{za sve } i$$

$$f(c_2) - L_2 d_2 \leq f(c_i) - L_2 d_i \quad \text{za sve } i$$

$$f(c_3) - L_2 d_3 \leq f(c_i) - L_2 d_i \quad \text{za sve } i$$



- svaki od navedenih pravaca određuje barem jednu potencijalno optimalnu točku;
- sve dobre točke leže na konveksnoj ljusci;

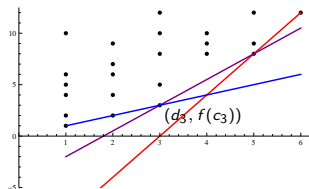
Traženje potencijalno optimalnih intervala

Međutim, kao što se može vidjeti ima i boljih pravaca s većim koeficijentom smjera, a koji imaju spomenuto svojstvo:

$$f(c_1) - L_1 d_1 \leq f(c_i) - L_1 d_i \quad \text{za sve } i$$

$$f(c_2) - L_2 d_2 \leq f(c_i) - L_2 d_i \quad \text{za sve } i$$

$$f(c_3) - L_3 d_3 \leq f(c_i) - L_3 d_i \quad \text{za sve } i$$



- svaki od navedenih pravaca određuje barem jednu potencijalno optimalnu točku;
- sve dobre točke leže na konveksnoj ljusci;

Zašto nije dovoljno dobro izabrati točke s najmanjom vrijednosti funkcije ?

$$f: [1, 7] \rightarrow \mathbb{R}, \quad f(x) = \begin{cases} 2(x-1)^2 + 3, & 0 \leq x \leq 2, \\ (x-4)^2 + 1, & 2 \leq x \leq 5, \\ 2, & 5 \leq x \leq 5.5, \\ (x-6.5)^2 + 1, & 5.5 \leq x \leq 7 \end{cases}$$

- $m = 30$; $a = 0$; $b = 10$.; $y = \text{RandomReal}[\{a, b\}, m]$;
- $f: [0, 10] \rightarrow [0, 20]$, $f(x) = \underset{i=1, \dots, m}{\text{med}} (y_i - x)^2$;
- $L = 9$;

Lema (Gablonsky,2001)

Lemma 2

Neka je $f \in Lip_L[a, b]$ Lipschitz neprekidna funkcija s konstantom $L > 0$, $\{d_i, c_i : i \in I\}$ poluširine i centri intervala nastali dijeljenjem intervala $[a, b]$ u algoritmu DIRECT, $f_{min} = \min_{i \in I} f(c_i)$ trenutna vrijednost minimuma.

Interval $[a_j, b_j]$, $j \in I$ je potencijalno optimalan onda i samo onda ako je

$$f(c_j) \leq f(c_i), \quad \forall i \in I_0, \quad (1)$$

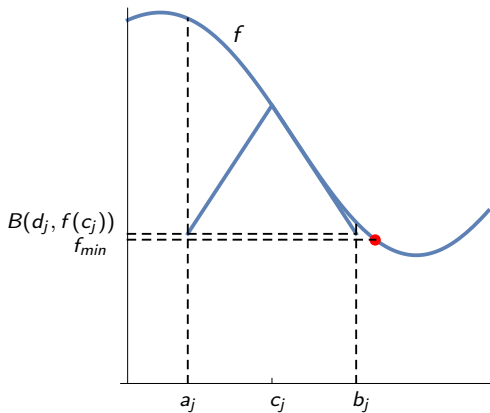
$$\max_{k \in I_L} \frac{f(c_j) - f(c_k)}{d_j - d_k} \leq \min_{k \in I_R} \frac{f(c_j) - f(c_k)}{d_j - d_k}, \quad (2)$$

gdje je

- $I_0 = \{i \in I : d_i = d_j\}$ – indeksi točaka s apscisom d_j ;
- $I_L = \{i \in I : d_i < d_j\}$ – indeksi točaka koje se na grafikonu nalaze lijevo od točaka s apscisom d_j ;
- $I_R = \{i \in I : d_i > d_j\}$ – indeksi točaka koje se na grafikonu nalaze desno od točaka s apscisom d_j .

Ako je $[a_j, b_j]$ potencijalno optimalan intrval sukladno, onda vrijedi

$$\mathcal{B}(d_j, f(c_j)) = f(c_j) - \hat{L}d_j \leq f(c_j) \leq f_{min}.$$



Slika: Interval koji nije potencijalno optimalan

Lemma 3

Neka je $f \in Lip_L[a, b]$ Lipschitz neprekidna funkcija s konstantom $L > 0$, $\epsilon > 0$, $\{d_i, c_i : i \in I = \{1, \dots, n\}$, poluširine i centri intervala nastali dijeljenjem intervala $[a, b]$ u algoritmu DIRECT i $f_{min} = \min_{i \in I} f(c_i)$.

Neka je $S \subset I$ skup indeksa potencijalno optimalnih intervala i neka je $\epsilon > 0$. B -vrijednost potencijalno optimalnog intervala $[a_j, b_j]$, $j \in S$ razlikuje se od aktualnog minimuma f_{min} za više od ϵ onda i samo onda ako vrijedi

$$\epsilon \leq \frac{f_{min} - f(c_j)}{|f_{min}|} + \frac{d_j}{|f_{min}|} \min_{k \in I_R} \frac{f(c_j) - f(c_k)}{d_j - d_k}, \quad \text{ako je } f_{min} \neq 0, \quad (3)$$

odnosno

$$f(c_j) \leq d_j \min_{k \in I_R} \frac{f(c_j) - f(c_k)}{d_j - d_k}, \quad \text{ako je } f_{min} = 0, \quad (4)$$

gdje je $I_R = \{i \in I : d_i > d_j\}$.

Theorem 2

Neka je $f \in Lip_L[a, b]$ Lipschitz neprekidna funkcija s konstantom $L > 0$, $\{d_i, c_i : i \in I\}$ poluširine i centri intervala nastali dijeljenjem intervala $[a, b]$ u algoritmu DIRECT, $f_{min} = \min_{i \in I} f(c_i)$ aktualna vrijednost minimuma $i \in I$.

Interval $[a_j, b_j]$, $j \in I$ je potencijalno optimalan, a njegova β -vrijednost razlikuje se od aktualnog minimuma f_{min} za više od ϵ onda i samo onda ako vrijedi

$$f(c_j) \leq f(c_i), \quad \forall i \in I_0,$$

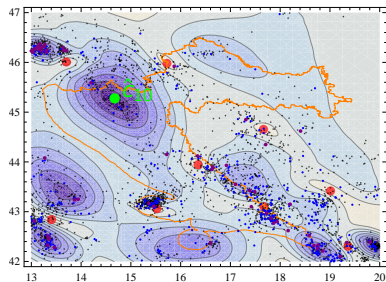
$$\max_{k \in I_L} \frac{f(c_j) - f(c_k)}{d_j - d_k} \leq \min_{k \in I_R} \frac{f(c_j) - f(c_k)}{d_j - d_k},$$

$$\begin{cases} \epsilon \leq \frac{f_{min} - f(c_j)}{|f_{min}|} + \frac{d_j}{|f_{min}|} \min_{k \in I_R} \frac{f(c_j) - f(c_k)}{d_j - d_k}, & \text{ako je } f_{min} \neq 0, \\ f(c_j) \leq d_j \min_{k \in I_R} \frac{f(c_j) - f(c_k)}{d_j - d_k}, & \text{ako je } f_{min} = 0, \end{cases}$$

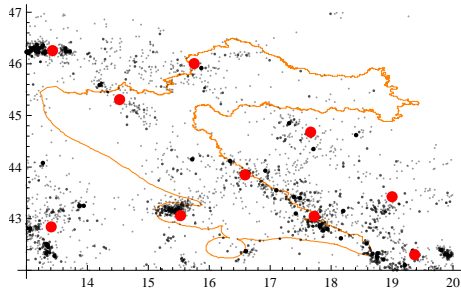
gdje je $I_0 = \{i \in I : d_i = d_j\}$, $I_L = \{i \in I : d_i < d_j\}$, $I_R = \{i \in I : d_i > d_j\}$.

Primjer: detekcija centara seizmičke aktivnosti

(a) DIRECT algoritam



(b) k-means algoritam



$$F: [13, 20] \times [42, 47] \rightarrow \mathbb{R}_+,$$

$$F(c) = \sum_{i=1}^m \min\{d(\hat{c}_1, a_i), \dots, d(\hat{c}_{k-1}, a_i), d(c, a_i)\}$$