The third International School on Model Reduction for Dynamical Control Systems October 5–10, 2015 Dubrovnik, Croatia

Model Order Reduction via System Balancing

Peter Benner

Max Planck Institute for Dynamics of Complex Technical Systems Computational Methods in Systems and Control Theory Magdeburg, Germany

benner@mpi-magdeburg.mpg.de

Outline



- Introduction
- Model Reduction for Dynamical Systems
- Application Areas
- Motivating Examples



- Mathematical Basics
- Numerical Linear Algebra
- Systems and Control Theory
- Qualitative and Quantitative Study of the Approximation Error
- Model Reduction by Projection
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 - Projection-based MOR Methods
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Outline



Introduction

- Model Reduction for Dynamical Systems
- Application Areas
- Motivating Examples

Mathematical Basics

Model Reduction by Projection









Introduction Model Reduction for Dynamical Systems

Dynamical Systems

$$\Sigma: \begin{cases} \dot{x}(t) = f(t, x(t), u(t)), & x(t_0) = x_0, \\ y(t) = g(t, x(t), u(t)) \end{cases}$$

with

• states
$$x(t) \in \mathbb{R}^n$$

• inputs
$$u(t) \in \mathbb{R}^m$$
,

• outputs $y(t) \in \mathbb{R}^q$.



Original System

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- states $\hat{x}(t) \in \mathbb{R}^r$, $r \ll n$
- inputs $u(t) \in \mathbb{R}^m$,
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$$\xrightarrow{u}$$
 $\hat{\Sigma}$ \hat{y}

Goal:

 $||y - \hat{y}|| < \text{tolerance} \cdot ||u||$ for all admissible input signals.

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• outputs
$$y(t) \in \mathbb{R}^{3}$$
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Goal:

 $\|y - \hat{y}\| < \text{tolerance} \cdot \|u\|$ for all admissible input signals.



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Goal:

 $||y - \hat{y}|| < \text{tolerance} \cdot ||u||$ for all admissible input signals. Secondary goal: reconstruct approximation of x from \hat{x} .

Linear, Time-Invariant (LTI) Systems

$$\begin{array}{rcl} \dot{x} & = & f(t,x,u) & = & Ax + Bu, & A \in \mathbb{R}^{n \times n}, \\ y & = & g(t,x,u) & = & Cx + Du, & C \in \mathbb{R}^{q \times n}, \\ \end{array} \qquad \begin{array}{rcl} B \in \mathbb{R}^{n \times m}, \\ D \in \mathbb{R}^{q \times m}. \end{array}$$

Introduction Mathematical Basics MOR by Projection Modal Truncation Balanced Truncation Matrix Equations

Application Areas Structural Mechanics / Finite Element Modeling

since ~ 1960 ies



- Resolving complex 3D geometries \Rightarrow millions of degrees of freedom.
- Analysis of elastic deformations requires many simulation runs for varying external forces, in particular if the model is used in an (elastic) multi-body simulation ((E)MBS).

Standard MOR techniques in structural mechanics: modal truncation, combined with Guyan reduction (static condensation) \rightsquigarrow Craig-Bampton method.

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- real-time constraints,
- increasing fragility for larger N.

 \implies reduce order of plant (*n*) and/or controller (*N*).



- input = output of plant,
- output = input of plant.

 $\begin{array}{l} \mbox{Modern (LQG-}/\mathcal{H}_{2^{-}}/\mathcal{H}_{\infty}\text{-}) \mbox{ control} \\ \mbox{design: } N \geq n. \end{array}$



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Progressive miniaturization

- Verification of VLSI/ULSI chip design requires high number of simulations for different input signals.
- Moore's Law (1965/75) states that the number of on-chip transistors doubles each 24 months.



Source: http://en.wikipedia.org/wiki/File:Transistor_Count_and_Moore'sLaw_-_2011.svg

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Intel 4004 (1971)	Intel Core 2 Extreme (quad-core) (2007)			
1 layer, 10μ technology	9 layers, 45 <i>nm</i> technology			
2,300 transistors	> 8,200,000 transistors			
64 kHz clock speed	> 3 GHz clock speed.			

since ${\sim}1990 \text{ies}$

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Source: http://en.wikipedia.org/wiki/Image:Silicon_chip_3d.png.

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- Here: mostly MOR for linear systems, they occur in micro electronics through modified nodal analysis (MNA) for RLC networks. e.g., when
 - decoupling large linear subcircuits,
 - modeling transmission lines,
 - modeling pin packages in VLSI chips,
 - modeling circuit elements described by Maxwell's equation using partial element equivalent circuits (PEEC).

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Standard MOR techniques in circuit simulation: Krylov subspace / Padé approximation / rational interpolation methods.

Application Areas

Many other disciplines in computational sciences and engineering like

- computational fluid dynamics (CFD),
- computational electromagnetics,
- chemical process engineering,
- design of MEMS/NEMS (micro/nano-electrical-mechanical systems),
- computational acoustics,
- . . .
- **Current trend:** more and more multi-physics problems, i.e., coupling of several field equations, e.g.,
 - electro-thermal (e.g., bondwire heating in chip design),
 - fluid-structure-interaction,
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Peter Benner and Lihong Feng.

Model Order Reduction for Coupled Problems

Applied and Computational Mathematics: An International Journal, 14(1):3-22, 2015. Available from http://www2.mpi-magdeburg.mpg.de/preprints/2015/MPIMD15-02.pdf

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- Main computational cost for set-up data $\approx 22 min$.
- Computation of reduced models from set-up data: 44–49sec. (r = 20-70).
- Bode plot (MATLAB on Intel Core i7, 2,67GHz, 12GB):
 7.5h for original system , < 1min for reduced system.



Motivating Examples Electro-Thermic Simulation of Integrated Circuit (IC)

[Source: Evgenii Rudnyi, CADFEM GmbH]

- Original model: n = 270.593, m = q = 2 ⇒
 Computing time (on Intel Xeon dualcore 3GHz, 1 Thread):
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Motivating Examples A Nonlinear Model from Computational Neurosciences: the FitzHugh-Nagumo System

• Simple model for neuron (de-)activation [Chaturantabut/Sorensen 2009]

$$\epsilon v_t(x,t) = \epsilon^2 v_{xx}(x,t) + f(v(x,t)) - w(x,t) + g$$

$$w_t(x,t) = hv(x,t) - \gamma w(x,t) + g,$$

with f(v) = v(v - 0.1)(1 - v) and initial and boundary conditions

$$egin{aligned} & v(x,0) = 0, & w(x,0) = 0, & x \in [0,1] \\ & v_x(0,t) = -i_0(t), & v_x(1,t) = 0, & t \geq 0, \end{aligned}$$

where $\epsilon = 0.015$, h = 0.5, $\gamma = 2$, g = 0.05, $i_0(t) = 50000t^3 \exp(-15t)$.



Source: http://en.wikipedia.org/wiki/Neuron

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- Parameter g handled as an additional input.
- Original state dimension $n = 2 \cdot 400$, QBDAE dimension $N = 3 \cdot 400$, reduced QBDAE dimension r = 26, chosen expansion point $\sigma = 1$.

Motivating Examples A Nonlinear Model from Computational Neurosciences: the FitzHugh-Nagumo System





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Mathematical Basics

- Numerical Linear Algebra
- Systems and Control Theory
- Qualitative and Quantitative Study of the Approximation Error

3 Model Reduction by Projection



Balanced Truncation





Numerical Linear Algebra Image Compression by Truncated SVD

- A digital image with $n_x \times n_y$ pixels can be represented as matrix $X \in \mathbb{R}^{n_x \times n_y}$, where x_{ji} contains color information of pixel (i, j).
- Memory (in single precision): $4 \cdot n_x \cdot n_y$ bytes.

Theorem (Schmidt-Mirsky/Eckart-Young)

Best rank-*r* approximation to $X \in \mathbb{R}^{n_x \times n_y}$ w.r.t. spectral norm:

$$\widehat{X} = \sum_{j=1}^r \sigma_j u_j v_j^T,$$

where $X = U\Sigma V^T$ is the singular value decomposition (SVD) of X. The approximation error is $||X - \hat{X}||_2 = \sigma_{r+1}$.

Idea for dimension reduction

Instead of X save $u_1, \ldots, u_r, \sigma_1 v_1, \ldots, \sigma_r v_r$. \rightsquigarrow memory = $4r \times (n_x + n_y)$ bytes.

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Example: Image Compression by Truncated SVD



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• rank r = 50, ≈ 104 kB


Event and the states which by rejection model transation balance transation water equations operations operati

Example: Image Compression by Truncated SVD



• rank r = 50, ≈ 104 kB



• rank r = 20, ≈ 42 kB

ank-20 approximation



Dimension Reduction via SVD

Example: Gatlinburg

Organizing committee Gatlinburg/Householder Meeting 1964: James H. Wilkinson, Wallace Givens, George Forsythe, Alston Householder, Peter Henrici, Fritz L. Bauer.



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• rank r = 100, ≈ 448 kB



• rank r = 50, ≈ 224 kB

Rank-50 approximation



Background: Singular Value Decay

Image data compression via SVD works, if the singular values decay (exponentially).



A different viewpoint

Linear Mapping

A matrix $A \in \mathbb{R}^{\ell imes k}$ represents a linear mapping

$$\mathcal{A}: \mathbb{R}^k \to \mathbb{R}^\ell : x \to y := Ax.$$

The truncated SVD ignores small singular values and thus the related left and right singular vectors.

Consequence:

- Vectors (almost) in the kernel of A do not contribute to range (A) and can hardly or not at all be reconstructed from the input-output relation (" A^{-1} ") \rightsquigarrow "unobservable" states.
- Vectors (almost) in range $(A)^{\perp}$ cannot be "reached" from any $x \in \mathbb{R}^k \rightsquigarrow$ "unreachable/uncontrollable" states.
- Hence, the truncated SVD ignores states hard to reconstruct and hard to reach.

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Systems and Control Theory The Laplace transform

Definition

The Laplace transform of a time domain function $f \in L_{1,\text{loc}}$ with $\text{dom}(f) = \mathbb{R}_0^+$ is

$$\mathcal{L}: f \mapsto F, \quad F(s) := \mathcal{L}{f(t)}(s) := \int_0^\infty e^{-st} f(t) dt, \quad s \in \mathbb{C}.$$

F is a function in the (Laplace or) frequency domain.

Note: for frequency domain evaluations ("frequency response analysis"), one takes re s = 0 and im $s \ge 0$. Then $\omega := \text{im } s$ takes the role of a frequency (in [rad/s], i.e., $\omega = 2\pi v$ with v measured in [Hz]).

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Lemma

$$\mathcal{L}\{\dot{f}(t)\}(s)=sF(s)-f(0).$$

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Note: for ease of notation, in the following we will use lower-case letters for both, a function and its Laplace transform!

Linear Systems in Frequency Domain

Application of Laplace transform $(x(t) \mapsto x(s), \dot{x}(t) \mapsto sx(s))$ to linear system

$$\dot{x}(t) = Ax(t) + Bu(t), \quad y(t) = Cx(t) + Du(t)$$

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 \Longrightarrow I/O-relation in frequency domain:

$$y(s) = \left(\underbrace{C(sI_n - A)^{-1}B + D}_{=:G(s)}\right)u(s).$$

G(s) is the transfer function of Σ .

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Goal: Fast evaluation of mapping $u \rightarrow y$.

Formulating model reduction in frequency domain

Approximate the dynamical system

$$\begin{aligned} \dot{x} &= Ax + Bu, \qquad A \in \mathbb{R}^{n \times n}, \ B \in \mathbb{R}^{n \times m}, \\ y &= Cx + Du, \qquad C \in \mathbb{R}^{q \times n}, \ D \in \mathbb{R}^{q \times m}, \end{aligned}$$

by reduced-order system

$$\begin{aligned} \dot{\hat{x}} &= \hat{A}\hat{x} + \hat{B}u, \quad \hat{A} \in \mathbb{R}^{r \times r}, \ \hat{B} \in \mathbb{R}^{r \times m}, \\ \hat{y} &= \hat{C}\hat{x} + \hat{D}u, \quad \hat{C} \in \mathbb{R}^{q \times r}, \ \hat{D} \in \mathbb{R}^{q \times m} \end{aligned}$$

of order $r \ll n$, such that

$$||y - \hat{y}|| = ||Gu - \hat{G}u|| \le ||G - \hat{G}|| \cdot ||u|| < \text{tolerance} \cdot ||u||$$

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by reduced-order system

$$\begin{aligned} \dot{\hat{x}} &= \hat{A}\hat{x} + \hat{B}u, \quad \hat{A} \in \mathbb{R}^{r \times r}, \ \hat{B} \in \mathbb{R}^{r \times m}, \\ \hat{y} &= \hat{C}\hat{x} + \hat{D}u, \quad \hat{C} \in \mathbb{R}^{q \times r}, \ \hat{D} \in \mathbb{R}^{q \times m} \end{aligned}$$

of order $r \ll n$, such that

$$\|y - \hat{y}\| = \|Gu - \hat{G}u\| \le \|G - \hat{G}\| \cdot \|u\| < \text{tolerance} \cdot \|u\|$$

 $\implies \text{Approximation problem: } \min_{\text{order}\,(\hat{G}) \leq r} \|G - \hat{G}\|.$

Definition

A linear system

$$\dot{x}(t) = Ax(t) + Bu(t), \quad y(t) = Cx(t) + Du(t)$$

is stable if its transfer function G(s) has all its poles in the left half plane and it is asymptotically (or Lyapunov or exponentially) stable if all poles are in the open left half plane $\mathbb{C}^- := \{z \in \mathbb{C} \mid \Re(z) < 0\}$.

Lemma

Sufficient for asymptotic stability is that A is asymptotically stable (or Hurwitz), i.e., the spectrum of A, denoted by $\Lambda(A)$, satisfies $\Lambda(A) \subset \mathbb{C}^-$.

Note that by abuse of notation, often *stable system* is used for asymptotically stable systems.

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Questions:

- For fixed $x_0 \in \mathbb{R}^n$ and some $x^1 \in \mathbb{R}^n$, is there a feasible control function $u \in \mathcal{U}_{ad}$ and time $t_1 > t_0 = 0$ such that $x(t_1; u) = x^1$? What is the set of targets x^1 reachable from x^0 ?
- For fixed $x_1 \in \mathbb{R}^n$ and some $x^0 \in \mathbb{R}^n$, is there a feasible control function $u \in \mathcal{U}_{ad}$ and time $t_1 > t_0 = 0$ such that $x(t_1; u) = x^1$? What is the set of initial conditions x^0 controllable to x^1 ?

Systems and Control Theory Properties of linear systems

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Systems and Control Theory Properties of linear systems

Questions:

- For fixed x₀ ∈ ℝⁿ and some x¹ ∈ ℝⁿ, is there a feasible control function u ∈ U_{ad} and time t₁ > t₀ = 0 such that x(t₁; u) = x¹? What is the set of targets x¹ reachable from x⁰?
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Note: for LTI systems $\dot{x} = Ax + Bu$, both concepts are equivalent!

Definition (Controllability)

Consider the target (the state to be reached) $x^1 \in \mathbb{R}^n$.

- a) An LTI system with initial value $x(0) = x^0$ is controllable to x^1 in time $t_1 > 0$ if there exists $u \in U_{ad}$ such that $x(t_1; u) = x^1$. (Equivalently, (t_1, x^1) is reachable from $(0, x^0)$.)
- b) x^0 is controllable to x^1 if there exists a $t_1 > 0$ such that (t_1, x^1) can be reached from $(0, x^0)$.
- c) If the system is controllable to x^1 for all $x^0 \in \mathbb{R}^n$, it is (completely) controllable.

The controllability set w.r.t. x^1 is defined as $\mathcal{C} := \bigcup_{t_1 > 0} \mathcal{C}(t_1)$ where

$$\mathcal{C}(t_1) := \{x^0 \in \mathbb{R}^n; \exists u \in \mathcal{U}_{ad} : x(t_1; u) = x^1\}.$$

In short: an LTI system is controllable $\iff C = \mathbb{R}^n$.

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Variation of constants \Longrightarrow

$$x(t) = e^{At}x^0 + \int_0^t e^{A(t-s)}Bu(s)ds = e^{At}(x^0 + \int_0^t e^{-As}Bu(s)ds).$$

Systems and Control Theory Properties of linear systems

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$$x^{1} = x(t_{1}) = e^{At_{1}}x^{0} + \int_{0}^{t_{1}} e^{A(t_{1}-t)}Bu(t)dt$$

This is equivalent to

$$e^{-At_1}x^1 - x^0 = \int_0^{t_1} e^{-At}Bu(t)dt.$$

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Ansatz: $u(t) = B^T e^{-A^T t} c \Longrightarrow$

$$e^{-At_1}x^1 - x^0 = \int_0^{t_1} e^{-At}BB^T e^{-A^T t}dtc =: P(0, t_1)c.$$

Systems and Control Theory Properties of linear systems

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Hence, an LTI system is controllable iff this linear system is solvable for $c \in \mathbb{R}^n$, i.e., iff $P(0, t_1)$ is invertible. (Note: $P(0, t_1) = P(0, t_1)^T \ge 0$ by definition!)

Now: characterize controllability.

Theorem

For an LTI system defined by $(A, B) \in \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times m}$, T.F.A.E.:

- a) The LTI system $\dot{x} = Ax + Bu$ is controllable.
- b) The finite time Gramian $P(0, t_1)$ is spd $\forall t_1 > 0$.
- c) The controllability matrix

$$K(A,B) := [B,AB,A^2B,\ldots,A^{n-1}B] \in \mathbb{R}^{n \times n \cdot m}$$

has full rank n. (Note: range $(K(A, B)) = C(t_1) \forall t_1 > 0!$)

- d) If z is a left eigenvector of A, then $z^*B \neq 0$.
- e) (Hautus test) rank($[\lambda I A, B]$) = $n \forall \lambda \in \mathbb{C}$.

Systems and Control Theory Properties of linear systems

The Gramian characterization of controllability for stable systems can be based on positive definiteness of the (infinite) controllability Gramian

$$\mathsf{P}:=\int_0^\infty e^{\mathsf{A}s}\mathsf{B}\mathsf{B}^\mathsf{T}e^{\mathsf{A}^\mathsf{T}s}\mathsf{d}s,$$

using congruence of $P(0, t_1)$ to $\int_{0}^{t_1} e^{As} BB^T e^{A^T s} ds$ and taking the limit $t_1 \to \infty$.

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using congruence of $P(0, t_1)$ to $\int_{0}^{t_1} e^{As}BB^T e^{A^Ts} ds$ and taking the limit $t_1 \to \infty$.

Theorem

For a stable LTI system defined by $(A, B) \in \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times m}$, T.F.A.E.:

- a) The LTI system $\dot{x} = Ax + Bu$ is controllable.
- b) The controllability Gramian P is positive definite.

New question: suppose we have

$$y(t) = \tilde{y}(t)$$

corresponding to two trajectories x, \tilde{x} obtained by the same input function u(t). Can we conclude that $x(0) = \tilde{x}(0)$, or even stronger, that $x(t) = \tilde{x}(t)$ for $t \le 0, t \ge 0$ (past/future)?

(Note that $x(t_0) = \tilde{x}(t_0)$ is sufficient as trajectory uniquely determined. In other words, is the mapping $x^0 \to y(t)$ injective?)

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Definition (Observability)

An LTI system is reconstructable (observable) if for solution trajectories $x(t), \tilde{x}(t)$ obtained with the same input function u, we have

$$y(t) = ilde y(t) \quad orall t \le 0 \quad (orall t \ge 0) \ \Longrightarrow \quad x(t) = \quad ilde x(t) \quad orall t \le 0 \quad (orall t \ge 0).$$

Systems and Control Theory Properties of linear systems

Characterization of observability/reconstructability:

Theorem (Duality)

An LTI system is reconstructable if and only if the dual system $\dot{x}(t) = -A^T x(t) - C^T u(t)$ is controllable.
Systems and Control Theory Properties of linear systems

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Theorem

For an LTI system defined by $(A, C) \in \mathbb{R}^{n \times n} \times \mathbb{R}^{q \times n}$, T.F.A.E.:

- a) The LTI system is reconstructable.
- b) The LTI system is observable.
- c) The observability matrix

$$\mathcal{O}(A,C) = \left[C^{T}, A^{T}C^{T}, (A^{2})^{T}C, \dots, (A^{n-1})^{T}C^{T}\right]^{T} \in \mathbb{R}^{np \times n} \text{ has rank } n.$$

d) If $Ax = \lambda x$, then $C^T x \neq 0$.

e) (Hautus test) rank
$$\begin{bmatrix} \lambda I - A \\ C \end{bmatrix} = n.$$

Systems and Control Theory Properties of linear systems

Characterization of observability/reconstructability:

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An LTI system is reconstructable if and only if the dual system $\dot{x}(t) = -A^T x(t) - C^T u(t)$ is controllable.

Theorem

A stable LTI system is observable if and only if the observability Gramian

$$Q := \int_{0}^{\infty} e^{A^{T}t} C^{T} C e^{At} dt$$

is symmetric positive definite.

- Controllability/observability are sometimes too strong.
- Weaker requirement: is there $u \in U_{ad}$ to steer x_0 to vicinity of x^1 ?
- For LTI systems, it suffices to consider $x^1 = 0!$
- Hence, is there $u \in U_{ad}$ so that $\lim_{t\to\infty} x(t; u) = 0$ $(\forall x^0 \in \mathbb{R}^n)$?
- If the answer is yes, then the LTI system is called stabilizable

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Theorem

For an LTI system defined by $(A, B) \in \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times m}$, T.F.A.E.:

a) The LTI system is stabilizable.

b)
$$\exists F \in \mathbb{R}^{m \times n}$$
 with $\Lambda(A + BF) \subset \mathbb{C}^{-}$.

- c) If $p^*A = \tilde{\lambda}p^*$ and $\operatorname{Re}(\lambda) \ge 0$, then $p^*B \neq 0$.
- d) $\operatorname{rank}([A \lambda I, B]) = n \quad \forall \lambda \in \mathbb{C} \text{ with } \operatorname{Re}(\lambda) \geq 0.$
- e) In the (controllability) Kalman decomposition of (A, B),

$$V^{\mathsf{T}} A V = \begin{bmatrix} A_1 & A_2 \\ 0 & A_3 \end{bmatrix}, V^{\mathsf{T}} B = \begin{bmatrix} B_1 \\ 0 \end{bmatrix},$$

we have $\Lambda(A_3) \subset \mathbb{C}^-$.

Systems and Control Theory Properties of linear systems

 \exists dual concept of stabilizability, analogous to duality of controllability and observability.

Definition (Detectability)

An LTI system is detectable if for any solution x(t) of $\dot{x} = Ax$ with $Cx(t) \equiv 0$ we have $\lim_{t\to\infty} x(t) = 0$. (We can not observe all of x, but the unobservable part is stable.)

Systems and Control Theory Properties of linear systems

 \exists dual concept of stabilizability, analogous to duality of controllability and observability.

Theorem

For an LTI system defined by $(A, C) \in \mathbb{R}^{n \times n} \times \mathbb{R}^{q \times n}$, T.F.A.E.:

- a) The LTI system is detectable.
- b) (A^{T}, C^{T}) is stabilizable.

c)
$$Ax = \lambda x$$
, $\operatorname{Re}(\lambda) \ge 0 \Rightarrow C^T x \neq 0$.

d) rank
$$\begin{bmatrix} \lambda I - A \\ C \end{bmatrix} = n$$
 for all λ , Re $(\lambda) \ge 0$.

e) In the observability Kalman decomposition of (A^{T}, C^{T}) ,

$$W^{\mathsf{T}}AW = \begin{bmatrix} A_1 & 0 \\ A_2 & A_3 \end{bmatrix}, CW = \begin{bmatrix} C_1 & 0 \end{bmatrix},$$

we have $\Lambda(A_3) \subset \mathbb{C}^-$.

Definition

For a linear (time-invariant) system

$$\Sigma: \begin{cases} \dot{x}(t) = Ax(t) + Bu(t), & \text{with transfer function} \\ y(t) = Cx(t) + Du(t), & G(s) = C(sI - A)^{-1}B + D, \end{cases}$$

the quadruple $(A, B, C, D) \in \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times m} \times \mathbb{R}^{q \times n} \times \mathbb{R}^{q \times m}$ is called a realization of Σ .

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Realizations are not unique!

Transfer function is invariant under state-space transformations,

$$\mathcal{T}: \left\{ \begin{array}{ccc} x & \rightarrow & Tx, \\ (A, B, C, D) & \rightarrow & (TAT^{-1}, TB, CT^{-1}, D) \end{array} \right.$$

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Realizations are not unique!

 $\label{eq:transfer} \begin{array}{l} \mbox{Transfer function is invariant under addition of uncontrollable/unobservable states:} \end{array}$

$$\frac{d}{dt} \begin{bmatrix} x \\ x_1 \end{bmatrix} = \begin{bmatrix} A & 0 \\ 0 & A_1 \end{bmatrix} \begin{bmatrix} x \\ x_1 \end{bmatrix} + \begin{bmatrix} B \\ B_1 \end{bmatrix} u(t), \quad y(t) = \begin{bmatrix} C & 0 \end{bmatrix} \begin{bmatrix} x \\ x_1 \end{bmatrix} + Du(t),$$
$$\frac{d}{dt} \begin{bmatrix} x \\ x_2 \end{bmatrix} = \begin{bmatrix} A & 0 \\ 0 & A_2 \end{bmatrix} \begin{bmatrix} x \\ x_2 \end{bmatrix} + \begin{bmatrix} B \\ 0 \end{bmatrix} u(t), \quad y(t) = \begin{bmatrix} C & C_2 \end{bmatrix} \begin{bmatrix} x \\ x_2 \end{bmatrix} + Du(t),$$

for arbitrary $A_j \in \mathbb{R}^{n_j \times n_j}$, j = 1, 2, $B_1 \in \mathbb{R}^{n_1 \times m}$, $C_2 \in \mathbb{R}^{q \times n_2}$ and any $n_1, n_2 \in \mathbb{N}$.

Definition

For a linear (time-invariant) system

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the quadruple $(A, B, C, D) \in \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times m} \times \mathbb{R}^{q \times n} \times \mathbb{R}^{q \times m}$ is called a realization of Σ .

Realizations are not unique!

Hence,

$$(A, B, C, D), \qquad \left(\begin{bmatrix} A & 0 \\ 0 & A_1 \end{bmatrix}, \begin{bmatrix} B \\ B_1 \end{bmatrix}, \begin{bmatrix} C & 0 \end{bmatrix}, D \right),$$
$$(TAT^{-1}, TB, CT^{-1}, D), \qquad \left(\begin{bmatrix} A & 0 \\ 0 & A_2 \end{bmatrix}, \begin{bmatrix} B \\ 0 \end{bmatrix}, \begin{bmatrix} C & C_2 \end{bmatrix}, D \right),$$

are all realizations of Σ !

Definition

For a linear (time-invariant) system

$$\Sigma: \begin{cases} \dot{x}(t) = Ax(t) + Bu(t), & \text{with transfer function} \\ y(t) = Cx(t) + Du(t), & G(s) = C(sI - A)^{-1}B + D, \end{cases}$$

the quadruple $(A, B, C, D) \in \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times m} \times \mathbb{R}^{q \times n} \times \mathbb{R}^{q \times m}$ is called a realization of Σ .

Definition

The McMillan degree of Σ is the unique minimal number $\hat{n} \ge 0$ of states necessary to describe the input-output behavior completely. A minimal realization is a realization $(\hat{A}, \hat{B}, \hat{C}, \hat{D})$ of Σ with order \hat{n} .

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Theorem

A realization (A, B, C, D) of a linear system is minimal \iff (A, B) is controllable and (A, C) is observable.

Definition

A realization (A, B, C, D) of a linear system Σ is balanced if its infinite controllability/observability Gramians P/Q satisfy

 $P = Q = \operatorname{diag} \{\sigma_1, \ldots, \sigma_n\} \quad (\text{w.l.o.g. } \sigma_j \geq \sigma_{j+1}, \ j = 1, \ldots, n-1).$

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When does a balanced realization exist? Assume A to be Hurwitz, i.e. $\Lambda(A) \subset \mathbb{C}^-$. Then:

Theorem

Given a stable minimal linear system Σ : (*A*, *B*, *C*, *D*), a balanced realization is obtained by the state-space transformation with

$$T_b := \Sigma^{-\frac{1}{2}} V^T R,$$

where $P = S^T S$, $Q = R^T R$ (e.g., Cholesky decompositions) and $SR^T = U\Sigma V^T$ is the SVD of SR^T .

Proof. Easy.

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 $\sigma_1, \ldots, \sigma_n$ are the Hankel singular values of Σ .

Note: $\sigma_1, \ldots, \sigma_n \ge 0$ as $P, Q \ge 0$ by definition, and $\sigma_1, \ldots, \sigma_n > 0$ in case of minimality!

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$$AP + PA^T + BB^T = 0, \quad A^TQ + QA + C^TC = 0.$$

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Proof. Exercise!

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The Hankel singular values (HSVs) of a stable minimal linear system are system invariants, i.e. they are unaltered by state-space transformations!

Systems and Control Theory Balanced Realizations

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The Hankel singular values (HSVs) of a stable minimal linear system are system invariants, i.e. they are unaltered by state-space transformations!

Proof. In balanced coordinates, the HSVs are $\Lambda(PQ)^{\frac{1}{2}}$. Now let

$$(\hat{A}, \hat{B}, \hat{C}, D) = (TAT^{-1}, TB, CT^{-1}, D)$$

be any transformed realization with associated controllability Lyapunov equation

$$0 = \hat{A}\hat{P} + \hat{P}\hat{A}^{T} + \hat{B}\hat{B}^{T} = TAT^{-1}\hat{P} + \hat{P}T^{-T}A^{T}T^{T} + TBB^{T}T^{T}.$$

This is equivalent to

$$0 = A(T^{-1}\hat{P}T^{-T}) + (T^{-1}\hat{P}T^{-T})A^{T} + BB^{T}.$$

The uniqueness of the solution of the Lyapunov equation implies that $\hat{P} = TPT^T$ and, analogously, $\hat{Q} = T^{-T}QT^{-1}$. Therefore,

$$\hat{P}\hat{Q} = TPQT^{-1}$$

showing that $\Lambda(\hat{P}\hat{Q}) = \Lambda(PQ) = \{\sigma_1^2, \dots, \sigma_n^2\}.$

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Remark

For non-minimal systems, the Gramians can also be transformed into diagonal matrices with the leading $\hat{n} \times \hat{n}$ submatrices equal to $\operatorname{diag}(\sigma_1, \ldots, \sigma_{\hat{n}})$, and

$$\hat{P}\hat{Q} = \operatorname{diag}(\sigma_1^2,\ldots,\sigma_{\hat{n}}^2,0,\ldots,0).$$

see [LAUB/HEATH/PAIGE/WARD 1987, TOMBS/POSTLETHWAITE 1987].

Qualitative and Quantitative Study of the Approximation Error System Norms

Consider transfer function

$$G(s) = C \left(sI - A \right)^{-1} B + D$$

and input functions $u\in \mathcal{L}_2^m\cong \mathcal{L}_2^m(-\infty,\infty)$, with the \mathcal{L}_2 -norm

$$\|u\|_2^2 := \frac{1}{2\pi} \int_{-\infty}^{\infty} u(j\omega)^H u(j\omega) \, d\omega.$$

Assume A (asympotically) stable: $\Lambda(A) \subset \mathbb{C}^- := \{z \in \mathbb{C} : re z < 0\}.$

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$$\int_{-\infty}^{\infty} y(j\omega)^{H} y(j\omega) \, d\omega \quad = \quad \int_{-\infty}^{\infty} u(j\omega)^{H} G(j\omega)^{H} G(j\omega) u(j\omega) \, d\omega$$

(Here, $\|.\|$ denotes the Euclidian vector or spectral matrix norm.)

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 $\implies y \in \mathcal{L}_2^q \cong \mathcal{L}_2^q(-\infty,\infty).$

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$$\|G\|_{\infty} := \sup_{\|u\|_{2} \neq 0} \frac{\|Gu\|_{2}}{\|u\|_{2}}$$

is well defined. It can be shown that

$$\|G\|_{\infty} = \sup_{\omega \in \mathbb{R}} \|G(j\omega)\| = \sup_{\omega \in \mathbb{R}} \sigma_{\max} \left(G(j\omega)\right).$$

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Sketch of proof: $\|G(\jmath\omega)u(\jmath\omega)\| \le \|G(\jmath\omega)\| \|u(\jmath\omega)\| \Rightarrow "\le "$. Construct u with $\|Gu\|_2 = \sup_{\omega \in \mathbb{R}} \|G(\jmath\omega)\| \|u\|_2$.

Qualitative and Quantitative Study of the Approximation Error System Norms

Consider transfer function

$$G(s) = C \left(sI - A \right)^{-1} B + D.$$

Hardy space \mathcal{H}_{∞}

Function space of matrix-/scalar-valued functions that are analytic and bounded in $\mathbb{C}^+.$

The \mathcal{H}_{∞} -norm is

$$\|F\|_{\infty} := \sup_{\mathrm{re}\,s>0} \sigma_{\max}\left(F(s)\right) = \sup_{\omega\in\mathbb{R}} \sigma_{\max}\left(F(j\omega)\right).$$

Stable transfer functions are in the Hardy spaces

- \mathcal{H}_{∞} in the SISO case (single-input, single-output, m=q=1);
- $\mathcal{H}^{q imes m}_{\infty}$ in the MIMO case (multi-input, multi-output, m > 1, q > 1).

Qualitative and Quantitative Study of the Approximation Error System Norms

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Consequence of Parseval identity/Plancherel Theorem

$$L_2(-\infty,\infty)\cong \mathcal{L}_2, \quad L_2(0,\infty)\cong \mathcal{H}_2$$

Consequently, 2-norms in time and frequency domains coincide!

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\mathcal{H}_∞ approximation error

Reduced-order model \Rightarrow transfer function $\hat{G}(s) = \hat{C}(sI_r - \hat{A})^{-1}\hat{B} + \hat{D}$. $\|y - \hat{y}\|_2 = \|Gu - \hat{G}u\|_2 \le \|G - \hat{G}\|_{\infty} \|u\|_2.$

 \implies compute reduced-order model such that $\|G - \hat{G}\|_{\infty} < tol!$ Note: error bound holds in time- and frequency domain due to Plancherel!

Qualitative and Quantitative Study of the Approximation Error System Norms

Consider stable transfer function

$$G(s) = C (sI - A)^{-1} B$$
, i.e. $D = 0$.

Hardy space \mathcal{H}_2

Function space of matrix-/scalar-valued functions that are analytic \mathbb{C}^+ and bounded w.r.t. the $\mathcal{H}_2\text{-norm}$

$$\begin{split} \|F\|_2 &:= \quad \frac{1}{2\pi} \left(\sup_{\operatorname{re} \sigma > 0} \int_{-\infty}^{\infty} \|F(\sigma + \jmath\omega)\|_F^2 \, d\omega \right)^{\frac{1}{2}} \\ &= \quad \frac{1}{2\pi} \left(\int_{-\infty}^{\infty} \|F(\jmath\omega)\|_F^2 \, d\omega \right)^{\frac{1}{2}}. \end{split}$$

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- \mathcal{H}_2 in the SISO case (single-input, single-output, m = q = 1);
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 $\begin{aligned} \mathcal{H}_2 \text{ approximation error for impulse response } (u(t) &= u_0 \delta(t)) \\ \text{Reduced-order model} \Rightarrow \text{transfer function } \hat{G}(s) &= \hat{C}(sI_r - \hat{A})^{-1}\hat{B}. \\ \|y - \hat{y}\|_2 &= \|Gu_0 \delta - \hat{G}u_0 \delta\|_2 \leq \|G - \hat{G}\|_2 \|u_0\|. \\ \Rightarrow \text{ compute reduced-order model such that } \|G - \hat{G}\|_2 < to!! \end{aligned}$

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Theorem (Practical Computation of the \mathcal{H}_2 -norm)

$$\|F\|_2^2 = \operatorname{tr}\left(B^T Q B\right) = \operatorname{tr}\left(C P C^T\right),$$

where P, Q are the controllability and observability Gramians of the corresponding LTI system.

Qualitative and Quantitative Study of the Approximation Error Approximation Problems

Output errors in time-domain

$$\begin{aligned} \|y - \hat{y}\|_2 &\leq \|G - \hat{G}\|_{\infty} \|u\|_2 &\Longrightarrow \|G - \hat{G}\|_{\infty} < \text{tol} \\ \|y - \hat{y}\|_{\infty} &\leq \|G - \hat{G}\|_2 \|u\|_2 &\Longrightarrow \|G - \hat{G}\|_2 < \text{tol} \end{aligned}$$

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\mathcal{H}_∞ -norm	best approximation problem for given reduced order r in
	general open; balanced truncation yields suboptimal solu-
	tion with computable \mathcal{H}_∞ -norm bound.
\mathcal{H}_2 -norm	necessary conditions for best approximation known; (local)
	optimizer computable with iterative rational Krylov algo-
	rithm (IRKA)
Hankel-norm	optimal Hankel norm approximation (AAK theory).
$\ G\ _H := \sigma_{\max}$	

Qualitative and Quantitative Study of the Approximation Error Computable error measures

Evaluating system norms is computationally very (sometimes too) expensive.

Other measures

- absolute errors $\|G(j\omega_j) \hat{G}(j\omega_j)\|_2$, $\|G(j\omega_j) \hat{G}(j\omega_j)\|_\infty$ $(j = 1, ..., N_\omega)$; • relative errors $\frac{\|G(j\omega_j) - \hat{G}(j\omega_j)\|_2}{\|G(j\omega_j)\|_2}$, $\frac{\|G(j\omega_j) - \hat{G}(j\omega_j)\|_\infty}{\|G(j\omega_j)\|_\infty}$;
- "eyeball norm", i.e. look at frequency response/Bode (magnitude) plot: for SISO system, log-log plot frequency vs. $|G(j\omega)|$ (or $|G(j\omega) - \hat{G}(j\omega)|$) in decibels, 1 dB $\simeq 20 \log_{10}(\text{value})$.

For MIMO systems, $q \times m$ array of plots G_{ij} .



Outline



Mathematical Basics



Model Reduction by Projection

- Introduction
- Projection-based MOR Methods



5 Balanced Truncation



7 Final Remarks

- Automatic generation of compact models.
- Satisfy desired error tolerance for all admissible input signals, i.e., want

 $||y - \hat{y}|| < \text{tolerance} \cdot ||u|| \qquad \forall u \in L_2(\mathbb{R}, \mathbb{R}^m).$

 \implies Need computable error bound/estimate!

- Preserve physical properties:
 - stability (poles of G in \mathbb{C}^-),
 - minimum phase (zeroes of G in \mathbb{C}^-),
 - passivity

 $\int_{-\infty}^{t} u(\tau)^{T} y(\tau) \, d\tau \ge 0 \quad \forall t \in \mathbb{R}, \quad \forall u \in L_{2}(\mathbb{R}, \mathbb{R}^{m}).$

Model Reduction by Projection Goals

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Model Reduction by Projection Projection Basics

Definition (Projector)

A projector is a matrix $P \in \mathbb{R}^{n \times n}$ with $P^2 = P$. Let $\mathcal{V} = \operatorname{range}(P)$, then P is projector onto \mathcal{V} . On the other hand, if $\{v_1, \ldots, v_r\}$ is a basis of \mathcal{V} and $V = [v_1, \ldots, v_r]$, then $P = V(V^T V)^{-1} V^T$ is a projector onto \mathcal{V} .

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Lemma (Projector Properties)

- If P = P^T, then P is an orthogonal projector (aka: Galerkin projection), otherwise an oblique projector (aka: Petrov-Galerkin projection).
- *P* is the identity operator on \mathcal{V} , i.e., $Pv = v \ \forall v \in \mathcal{V}$.
- I P is the complementary projector onto ker P.
- If \mathcal{V} is an A-invariant subspace corresponding to a subset of A's spectrum, then we call P a spectral projector.

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Model Reduction by Projection Projection-based MOR Methods

Methods:

- Modal Truncation
- Balanced Truncation
- Rational Interpolation (Padé-Approximation and (rational) Krylov Subspace Methods)
- many more...

Joint feature of these methods:

computation of reduced-order model (ROM) by projection!

Model Reduction by Projection Projection-based MOR Methods

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range
$$(V) = \mathcal{V}$$
, range $(W) = \mathcal{W}$, $W^T V = I_r$.

Then, with $\hat{x} = W^T x$, we obtain $x \approx V \hat{x}$ so that

$$\|x - \tilde{x}\| = \|x - V\hat{x}\|$$

and the reduced-order model is

$$\hat{A} := W^T A V, \quad \hat{B} := W^T B, \quad \hat{C} := C V, \quad (\hat{D} := D).$$

Projection-based MOR Methods

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Important observation:

• The state equation residual satisfies $\dot{\tilde{x}} - A\tilde{x} - Bu \perp W$, since

$$W^{T}\left(\dot{\tilde{x}} - A\tilde{x} - Bu\right) = W^{T}\left(VW^{T}\dot{x} - AVW^{T}x - Bu\right)$$

Model Reduction by Projection Model France France Matrix Equations

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Model Reduction by Projection Model Function Salanced Function Matrix Equations

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$$= \dot{\hat{x}} - \hat{A}\hat{x} - \hat{B}u = 0.$$

Outline



- Mathematical Basics
- 3 Model Reduction by Projection
 - Modal Truncation
 - Basic Principle
 - Dominant Pole Algorithm

5 Balanced Truncation

6 Solving Large-Scale Matrix Equations

7 Final Remarks

Modal Truncation

Basic method:

Assume A is diagonalizable, $T^{-1}AT = D_A$, project state-space onto A-invariant subspace $\mathcal{V} = \operatorname{span}(t_1, \ldots, t_r)$, $t_k = \operatorname{eigenvectors}$ corresp. to "dominant" modes / eigenvalues of A. Then with

 $V = T(:, 1:r) = [t_1, ..., t_r], \quad \tilde{W}^H = T^{-1}(1:r, :), \quad W = \tilde{W}(V^H \tilde{W})^{-1},$

reduced-order model is

$$\hat{A} := W^H A V = \operatorname{diag} \{\lambda_1, \dots, \lambda_r\}, \quad \hat{B} := W^H B, \quad \hat{C} = C V$$

Also computable by truncation:

$$T^{-1}AT = \begin{bmatrix} \hat{A} \\ A_2 \end{bmatrix}, \quad T^{-1}B = \begin{bmatrix} \hat{B} \\ B_2 \end{bmatrix}, \quad CT = [\hat{C}, C_2], \quad \hat{D} = D.$$

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Properties:

Simple computation for large-scale systems, using, e.g., Krylov subspace methods (Lanczos, Arnoldi), Jacobi-Davidson method.

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Properties:

Error bound:

$$\|G - \hat{G}\|_{\infty} \leq \|C_2\| \|B_2\| \frac{1}{\min_{\lambda \in \Lambda(A_2)} |\operatorname{Re}(\lambda)|}$$

Proof:

$$\begin{aligned} G(s) &= C(sl - A)^{-1}B + D = CTT^{-1}(sl - A)^{-1}TT^{-1}B + D \\ &= CT(sl - T^{-1}AT)^{-1}T^{-1}B + D \\ &= [\hat{C}, C_2] \begin{bmatrix} (sl_r - \hat{A})^{-1} \\ (sl_{n-r} - A_2)^{-1} \end{bmatrix} \begin{bmatrix} \hat{B} \\ B_2 \end{bmatrix} + D \\ &= \hat{G}(s) + C_2(sl_{n-r} - A_2)^{-1}B_2, \end{aligned}$$

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Proof:

$$G(s) = \hat{G}(s) + C_2(sI_{n-r} - A_2)^{-1}B_2,$$

observing that $\|G - \hat{G}\|_{\infty} = \sup_{\omega \in \mathbb{R}} \sigma_{\max}(C_2(\jmath \omega I_{n-r} - A_2)^{-1}B_2)$, and

$$C_2(\jmath\omega I_{n-r}-A_2)^{-1}B_2=C_2 {
m diag}\left(rac{1}{\jmath\omega-\lambda_{r+1}},\ldots,rac{1}{\jmath\omega-\lambda_n}
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Difficulties:

- Eigenvalues contain only limited system information.
- Dominance measures are difficult to compute. ([LITZ '79] use Jordan canoncial form; otherwise merely heuristic criteria, e.g., [VARGA '95]. Recent improvement: dominant pole algorithm.)
- Error bound not computable for really large-scale problems.



BEAM, SISO system from SLICOT Benchmark Collection for Model Reduction, n = 348, m = q = 1, reduced using 13 dominant complex conjugate eigenpairs, error bound yields $||G - \hat{G}||_{\infty} \le 1.21 \cdot 10^3$



MATLAB[®] demo.



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MATLAB[®] demo.

Base enrichment

Static modes are defined by setting $\dot{x} = 0$ and assuming unit loads, i.e., $u(t) \equiv e_j, j = 1, ..., m$:

$$0 = Ax(t) + Be_j \implies x(t) \equiv -A^{-1}b_j.$$

Projection subspace \mathcal{V} is then augmented by $A^{-1}[b_1, \dots, b_m] = A^{-1}B$. Interpolation-projection framework $\implies G(0) = \hat{G}(0)!$

If two sided projection is used, complimentary subspace can be augmented by $A^{-T}C^T \Longrightarrow G'(0) = \hat{G}'(0)!$

Note: if $m \neq q$, add random vectors or delete some of the columns in $A^{-T}C^{T}$.

Guyan reduction (static condensation)

Partition states in masters $x_1 \in \mathbb{R}^r$ and slaves $x_2 \in \mathbb{R}^{n-r}$ (FEM terminology) Assume stationarity, i.e., $\dot{x} = 0$ and solve for x_2 in

$$0 = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} B_1 \\ B_2 \end{bmatrix} u$$

$$\Rightarrow \quad x_2 = -A_{22}^{-1}A_{21}x_1 - A_{22}^{-1}B_2u.$$

Inserting this into the first part of the dynamic system

$$\dot{x}_1 = A_{11}x_1 + A_{12}x_2 + B_1u, \quad y = C_1x_1 + C_2x_2$$

then yields the reduced-order model

=

$$\dot{x}_1 = (A_{11} - A_{12}A_{22}^{-1}A_{21})x_1 + (B_1 - A_{12}A_{22}^{-1}B_2)u y = (C_1 - C_2A_{22}^{-1}A_{21})x_1 - C_2A_{22}^{-1}B_2u.$$

Modal Truncation Dominant Pole Algorithm

Pole-Residue Form of Transfer Function

Consider partial fraction expansion of transfer function with D = 0:

$$G(s) = \sum_{k=1}^{n} \frac{R_k}{s - \lambda_k}$$

with the residues $R_k := (Cx_k)(y_k^H B) \in \mathbb{C}^{q \times m}$.

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Note: this follows using the spectral decomposition $A = XDX^{-1}$, with $X = [x_1, ..., x_n]$ the right and $X^{-1} =: Y = [y_1, ..., y_n]^H$ the left eigenvector matrices:

$$G(s) = C(sI - XDX^{-1})^{-1}B = CX(sI - \operatorname{diag} \{\lambda_1, \dots, \lambda_n\})^{-1}YB$$

$$= [Cx_1, \dots, Cx_n] \begin{bmatrix} \frac{1}{s-\lambda_1} & & \\ & \ddots & \\ & & \frac{1}{s-\lambda_n} \end{bmatrix} \begin{bmatrix} y_1^HB \\ \vdots \\ y_n^HB \end{bmatrix}$$

$$= \sum_{k=1}^n \frac{(Cx_k)(y_k^HB)}{s-\lambda_k}.$$

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Note: $R_k = (Cx_k)(y_k^H B)$ are the residues of G in the sense of the residue theorem of complex analysis:

$$\operatorname{res} (G, \lambda_{\ell}) = \lim_{s \to \lambda_{\ell}} (s - \lambda_{\ell}) G(s) = \sum_{k=1}^{n} \underbrace{\lim_{s \to \lambda_{\ell}} \frac{s - \lambda_{\ell}}{s - \lambda_{k}}}_{= \begin{cases} 0 \text{ for } k \neq \ell \\ 1 \text{ for } k = \ell \end{cases}} R_{k} = R_{\ell}.$$
Introduction
 Mathematical Basics
 MOR by Projection
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As projection basis use spaces spanned by right/left eigenvectors corresponding to dominant poles, i.e., (λ_i, x_i, y_i) with largest

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Remark

The dominant modes have most important influence on the input-output behavior of the system and are responsible for the "peaks"' in the frequency response.

















Outline



- 2 Mathematical Basics
- 3 Model Reduction by Projection

4 Modal Truncation



Balanced Truncation

- The basic method
- Theoretical Background
- Singular Perturbation Approximation
- Balancing-Related Methods





Basic principle:

 Recall: a stable system Σ, realized by (A, B, C, D), is called balanced, if the Gramians, i.e., solutions P, Q of the Lyapunov equations

$$AP + PA^T + BB^T = 0, \qquad A^TQ + QA + C^TC = 0,$$

satisfy: $P = Q = \operatorname{diag}(\sigma_1, \ldots, \sigma_n)$ with $\sigma_1 \ge \sigma_2 \ge \ldots \ge \sigma_n > 0$.

• $\Lambda(PQ)^{\frac{1}{2}} = \{\sigma_1, \dots, \sigma_n\}$ are the Hankel singular values (HSVs) of Σ .

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- Compute balanced realization of the system via state-space transformation

$$\mathcal{T} : (A, B, C, D) \quad \mapsto \quad (TAT^{-1}, TB, CT^{-1}, D) \\ = \quad \left(\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}, \begin{bmatrix} B_1 \\ B_2 \end{bmatrix}, \begin{bmatrix} C_1 & C_2 \end{bmatrix}, D \right)$$

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Motivation:

The HSVs $\Lambda(PQ)^{\frac{1}{2}} = \{\sigma_1, \dots, \sigma_n\}$ are system invariants: they are preserved under

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in transformed coordinates, the Gramians satisfy

$$(TAT^{-1})(TPT^{T}) + (TPT^{T})(TAT^{-1})^{T} + (TB)(TB)^{T} = 0,$$

$$(TAT^{-1})^{T}(T^{-T}QT^{-1}) + (T^{-T}QT^{-1})(TAT^{-1}) + (CT^{-1})^{T}(CT^{-1}) = 0$$

$$\Rightarrow (TPT^{T})(T^{-T}QT^{-1}) = TPQT^{-1},$$

hence $\Lambda(PQ) = \Lambda((TPT^{T})(T^{-T}QT^{-1})).$

Implementation: SR Method

• Compute (Cholesky) factors of the Gramians, $P = S^T S$, $Q = R^T R$.

- Compute SVD $SR^T = \begin{bmatrix} U_1, U_2 \end{bmatrix} \begin{bmatrix} \Sigma_1 \\ \Sigma_2 \end{bmatrix} \begin{bmatrix} V_1^T \\ V_2^T \end{bmatrix}.$
- Solution ROM is $(W^T AV, W^T B, CV, D)$, where

$$W = R^T V_1 \Sigma_1^{-\frac{1}{2}}, \qquad V = S^T U_1 \Sigma_1^{-\frac{1}{2}}$$

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 $\implies VW^{T}$ is an oblique projector, hence balanced truncation is a Petrov-Galerkin projection method.

Properties:

- Reduced-order model is stable with HSVs $\sigma_1, \ldots, \sigma_r$.
- Adaptive choice of *r* via computable error bound:

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Assumptions (for now): $t_0 = 0$, $x_0 = x(0) = 0$, D = 0.

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State-Space Description for $\mathsf{I}/\mathsf{O}\text{-}\mathsf{Relation}$

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Instead of

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Introduction Mathematical Basics MOR by Projection Modal Truncation Balanced Truncation Matrix Equations

Balanced Truncation Theoretical Background

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10²⁰ -HSVs \mathcal{H} compact machine precision 10 \mathcal{H} has discrete SVD p[≠] 10⁻²⁰ 11 10-46 Hankel singular values 10-60 100 200 300 400 500

Hankel Singular Values for Atmospheric Storm Model

600

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Let P, Q be the controllability and observability Gramians of an LTI system Σ . Then the Hankel singular values $\Lambda(PQ)^{\frac{1}{2}} = \{\sigma_1, \ldots, \sigma_n\}$ are the singular values of the Hankel operator associated to Σ .

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Theorem

Let the reduced-order system $\hat{\Sigma} : (\hat{A}, \hat{B}, \hat{C}, \hat{D})$ with $r \leq \hat{n}$ be computed by balanced truncation. Then the reduced-order model $\hat{\Sigma}$ is balanced, stable, minimal, and its HSVs are $\sigma_1, \ldots, \sigma_r$.

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Proof: Note that in balanced coordinates, the Gramians are diagonal and equal to

$$\operatorname{diag}(\boldsymbol{\Sigma}_1,\boldsymbol{\Sigma}_2)=\operatorname{diag}(\sigma_1,\ldots,\sigma_r,\sigma_{r+1},\ldots,\sigma_n).$$

Hence, the Gramian satisfies

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \begin{bmatrix} \Sigma_1 \\ \Sigma_2 \end{bmatrix} + \begin{bmatrix} \Sigma_1 \\ \Sigma_2 \end{bmatrix} \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}^T + \begin{bmatrix} B_1 \\ B_2 \end{bmatrix} \begin{bmatrix} B_1 \\ B_2 \end{bmatrix}^T = 0,$$

whence we obtain the "controllability Lyapunov equation" of the reduced-order system,

$$A_{11}\Sigma_1 + \Sigma_1 A_{11}^T + B_1 B_1^T = 0.$$

The result follows from $\hat{A} = A_{11}$, $\hat{B} = B_1$, $\Sigma_1 > 0$, the solution theory of Lyapunov equations and the analogous considerations for the observability Gramian. (Minimality is a simple consequence of $\hat{P} = \Sigma_1 = \hat{Q} > 0$.)

Assume the system

$$\begin{bmatrix} \dot{x}_1 \\ \dot{x}_2 \end{bmatrix} = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} B_1 \\ B_2 \end{bmatrix} u, \quad y = \begin{bmatrix} C_1, \ C_2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + Du$$

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$$\dot{x}_1 = (A_{11} - A_{12}A_{22}^{-1}A_{21})x_1 + (B_1 - A_{12}A_{22}^{-1}B_2)u, y = (C_1 - C_2A_{22}^{-1}A_{21})x_1 + (D - C_2A_{22}^{-1}B_2)u,$$

with

- the same properties as the reduced-order model w.r.t. stability, minimality, error bound, but $\hat{D} \neq D$;
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Note:

- A_{22} invertible as in balanced coordinates, $A_{22}\Sigma_2 + \Sigma_2 A_{22}^T + B_2 B_2^T = 0$ and (A_{22}, B_2) controllable, $\Sigma_2 > 0 \Rightarrow A_{22}$ stable.
- If the original system is not balanced, first compute a minimal realization by applying balanced truncation with $r = \hat{n}$.

Balancing-Related Methods

Basic Principle

Given positive semidefinite matrices $P = S^T S$, $Q = R^T R$, compute balancing state-space transformation so that

$$P = Q = \operatorname{diag}(\sigma_1, \ldots, \sigma_n) = \Sigma, \quad \sigma_1 \ge \ldots \ge \sigma_n > 0,$$

and truncate corresponding realization at size r with $\sigma_r > \sigma_{r+1}$.

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Classical Balanced Truncation (BT)

Mullis/Roberts '76, Moore '81

- P =controllability Gramian of system given by (A, B, C, D).
- Q = observability Gramian of system given by (A, B, C, D).
- P, Q solve dual Lyapunov equations

$$AP + PA^{T} + BB^{T} = 0, \qquad A^{T}Q + QA + C^{T}C = 0.$$

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LQG Balanced Truncation (LQGBT)

- P/Q = controllability/observability Gramian of closed-loop system based on LQG compensator.
- P, Q solve dual algebraic Riccati equations (AREs)

$$0 = AP + PA^{T} - PC^{T}CP + B^{T}B,$$

$$0 = A^{T}Q + QA - QBB^{T}Q + C^{T}C.$$

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Given positive semidefinite matrices $P = S^T S$, $Q = R^T R$, compute balancing state-space transformation so that

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Balanced Stochastic Truncation (BST) [Desai/Pal '84, Green '88]

- $P = \text{controllability Gramian of system given by } (A, B, C, D), \text{ i.e., solution of Lyapunov equation } AP + PA^T + BB^T = 0.$
- Q = observability Gramian of right spectral factor of power spectrum of system given by (A, B, C, D), i.e., solution of ARE

$$\hat{A}^{\mathsf{T}}Q + Q\hat{A} + QB_{W}(DD^{\mathsf{T}})^{-1}B_{W}^{\mathsf{T}}Q + C^{\mathsf{T}}(DD^{\mathsf{T}})^{-1}C = 0,$$

where $\hat{A} := A - B_W (DD^T)^{-1}C$, $B_W := BD^T + PC^T$.

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Positive-Real Balanced Truncation (PRBT)

- Based on positive-real equations, related to positive real (Kalman-Yakubovich-Popov-Anderson) lemma.
- P, Q solve dual AREs

$$0 = \bar{A}P + P\bar{A}^{T} + PC^{T}\bar{R}^{-1}CP + B\bar{R}^{-1}B^{T},$$

$$0 = \bar{A}^{T}Q + Q\bar{A} + QB\bar{R}^{-1}B^{T}Q + C^{T}\bar{R}^{-1}C.$$

where $\bar{R} = D + D^T$, $\bar{A} = A - B\bar{R}^{-1}C$.

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Balancing-Related Methods

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Other Balancing-Based Methods

- Bounded-real balanced truncation (BRBT) based on bounded real lemma [OPDENACKER/JONCKHEERE '88];
- H_{∞} balanced truncation (HinfBT) closed-loop balancing based on H_{∞} compensator [MUSTAFA/GLOVER '91].

Both approaches require solution of dual AREs.

• Frequency-weighted versions of the above approaches.

Balancing-Related Methods Properties

- Guaranteed preservation of physical properties like
 - stability (all),
 - passivity (PRBT),
 - minimum phase (BST).
- Computable error bounds, e.g.,

$$\begin{split} \mathsf{BT:} & \|G - G_r\|_{\infty} &\leq 2 \; \sum_{j=r+1}^{n} \sigma_j^{BT}, \\ \mathsf{LQGBT:} & \|G - G_r\|_{\infty} &\leq \; 2 \sum_{j=r+1}^{n} \frac{\sigma_j^{LQG}}{\sqrt{1 + (\sigma_j^{LQG})^2}} \\ \mathsf{BST:} & \|G - G_r\|_{\infty} &\leq \left(\prod_{j=r+1}^{n} \frac{1 + \sigma_j^{BST}}{1 - \sigma_j^{BST}} - 1\right) \|G\|_{\infty}, \end{split}$$

- Can be combined with singular perturbation approximation for steady-state performance.
- Computations can be modularized.

Outline

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- 2 Mathematical Basics
- 3 Model Reduction by Projection
- 4 Modal Truncation

5 Balanced Truncation

6 Solving Large-Scale Matrix Equations

- Linear Matrix Equations
- Numerical Methods for Solving Lyapunov Equations
- Solving Large-Scale Algebraic Riccati Equations
- Software

🕜 Final Remarks

Solving Large-Scale Matrix Equations Large-Scale Algebraic Lyapunov and Riccati Equations

Algebraic Riccati equation (ARE) for $A, G = G^T, W = W^T \in \mathbb{R}^{n \times n}$ given and $X \in \mathbb{R}^{n \times n}$ unknown:

 $0 = \mathcal{R}(X) := A^T X + X A - X G X + W.$

 $G = 0 \Longrightarrow$ Lyapunov equation:

 $0 = \mathcal{L}(X) := A^T X + X A + W.$

Typical situation in model reduction and optimal control problems for semi-discretized PDEs:

- $n = 10^3 10^6 \iff 10^6 10^{12}$ unknowns!),
- A has sparse representation $(A = -M^{-1}S$ for FEM),
- G, W low-rank with $G, W \in \{BB^T, C^T C\}$, where $B \in \mathbb{R}^{n \times m}$, $m \ll n$, $C \in \mathbb{R}^{p \times n}$, $p \ll n$.
- Standard (eigenproblem-based) O(n³) methods are not applicable!
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Consider spectrum of ARE solution (analogous for Lyapunov equations).

Example:

- Linear 1D heat equation with point control,
- Ω = [0, 1],

Low-Rank Approximation

- FEM discretization using linear B-splines,
- $h = 1/100 \implies n = 101.$

Idea: $X = X^T \ge 0 \implies$



$$X = ZZ^{T} = \sum_{k=1}^{n} \lambda_{k} z_{k} z_{k}^{T} \approx Z^{(r)} (Z^{(r)})^{T} = \sum_{k=1}^{r} \lambda_{k} z_{k} z_{k}^{T}.$$

 \implies Goal: compute $Z^{(r)} \in \mathbb{R}^{n \times r}$ directly w/o ever forming X!

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Equations without symmetry

Sylvester equation discrete Sylvester equation

AX + XB = W AXB - X = W

with data $A \in \mathbb{R}^{n \times n}$, $B \in \mathbb{R}^{m \times m}$, $W \in \mathbb{R}^{n \times m}$ and unknown $X \in \mathbb{R}^{n \times m}$.

Equations with symmetry

Lyapunov equation

Stein equation (discrete Lyapunov equation)

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 $AXA^{T} - X = W$

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Here: focus on (Sylvester and) Lyapunov equations; analogous results and methods for discrete versions exist.

Max Planck Institute Magdeburg

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Max Planck Institute Magdeburg

Using the Kronecker (tensor) product, AX + XB = W is equivalent to

$$((I_m \otimes A) + (B^T \otimes I_n)) \operatorname{vec} (X) = \operatorname{vec} (W).$$

Hence,

Sylvester equation has a unique solution

 $M := (I_m \otimes A) + (B^T \otimes I_n) \text{ is invertible.}$ \Leftrightarrow $0 \notin \Lambda(M) = \Lambda((I_m \otimes A) + (B^T \otimes I_n)) = \{\lambda_j + \mu_k \mid \lambda_j \in \Lambda(A), \ \mu_k \in \Lambda(B)\}.$ \Leftrightarrow $\Lambda(A) \cap \Lambda(-B) = \emptyset$

Corollary

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Linear Matrix Equations Solvability

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Corollary

Linear Matrix Equations Complexity Issues

Solving the Sylvester equation

$$AX + XB = W$$

via the equivalent linear system of equations

$$((I_m \otimes A) + (B^T \otimes I_n)) \operatorname{vec} (X) = \operatorname{vec} (W)$$

requires

- LU factorization of $nm \times nm$ matrix; for $n \approx m$, complexity is $\frac{2}{3}n^6$;
- storing $n \cdot m$ unknowns: for $n \approx m$ we have n^2 data for X, but up to n^4 data for triangular factors!

Example

 $n = m = 1,000 \Rightarrow$ Gaussian elimination on an Intel core i7 (Westmere, 6 cores, 3.46 GHz \rightarrow 83.2 GFLOP peak) would take > 94 DAYS and 7.3 TB of memory!

Linear Matrix Equations Complexity Issues

Solving the Sylvester equation

$$AX + XB = W$$

via the equivalent linear system of equations

$$((I_m \otimes A) + (B^T \otimes I_n)) \operatorname{vec} (X) = \operatorname{vec} (W)$$

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- LU factorization of $nm \times nm$ matrix; for $n \approx m$, complexity is $\frac{2}{3}n^6$;
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Numerical Methods for Solving Lyapunov Equations Traditional Methods

Bartels-Stewart method for Sylvester and Lyapunov equation (1yap); Hessenberg-Schur method for Sylvester equations (1yap); Hammarling's method for Lyapunov equations $AX + XA^T + GG^T = 0$ with A Hurwitz (1yapchol).

All based on the fact that if A, B^T are in Schur form, then

 $M = (I_m \otimes A) + (B^T \otimes I_n)$

- is block-upper triangular. Hence, solve Mx = b by back-substitution.
 - Clever implementation of back-substitution process requires nm(n+m) flops.
 - For Sylvester equations, *B* in Hessenberg form is enough (~~ Hessenberg-Schur method).
 - Hammarling's method computes Cholesky factor Y of X directly.
 - All methods require Schur decomposition of A and Schur or Hessenberg decomposition of $B \Rightarrow$ need QR algorithm which requires $25n^3$ flops for Schur decomposition.

Not feasible for large-scale problems (n > 10,000).

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Numerical Methods for Solving Lyapunov Equations The Sign Function Method

Definition

For $Z \in \mathbb{R}^{n \times n}$ with $\Lambda(Z) \cap i\mathbb{R} = \emptyset$ and Jordan canonical form

$$Z = S \begin{bmatrix} J^+ & 0 \\ 0 & J^- \end{bmatrix} S^{-1}$$

the matrix sign function is

$$\operatorname{sign}(Z) := S \left[\begin{array}{cc} I_k & 0 \\ 0 & -I_{n-k} \end{array} \right] S^{-1}.$$

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Lemma

Let $T \in \mathbb{R}^{n \times n}$ be nonsingular and Z as before, then

$$\operatorname{sign}\left(TZT^{-1}\right) = T\operatorname{sign}\left(Z\right)T^{-1}$$

Max Planck Institute Magdeburg

Numerical Methods for Solving Lyapunov Equations The Sign Function Method

Computation of sign (Z)

 $\operatorname{sign}(Z)$ is root of $I_n \Longrightarrow$ use Newton's method to compute it:

$$Z_0 \leftarrow Z, \qquad Z_{j+1} \leftarrow \frac{1}{2} \left(c_j Z_j + \frac{1}{c_j} Z_j^{-1} \right), \qquad j = 1, 2, \dots$$

$$\implies \operatorname{sign}(Z) = \lim_{j \to \infty} Z_j.$$

 $c_{\rm j}>0$ is scaling parameter for convergence acceleration and rounding error minimization, e.g.

$$c_j = \sqrt{\frac{\|Z_j^{-1}\|_F}{\|Z_j\|_F}},$$

based on "equilibrating" the norms of the two summands [HIGHAM '86].

Key observation:

If $X \in \mathbb{R}^{n \times n}$ is a solution of $AX + XA^T + W = 0$, then

$$\underbrace{\begin{bmatrix} I_n & -X \\ 0 & I_n \end{bmatrix}}_{=T^{-1}} \underbrace{\begin{bmatrix} A & W \\ 0 & -A^T \end{bmatrix}}_{=:H} \underbrace{\begin{bmatrix} I_n & X \\ 0 & I_n \end{bmatrix}}_{=:T} = \begin{bmatrix} A & 0 \\ 0 & -A^T \end{bmatrix}.$$

Hence, if A is Hurwitz (i.e., asymptotically stable), then

$$\operatorname{sign}(H) = \operatorname{sign}\left(T\begin{bmatrix}A & 0\\ 0 & -A^{T}\end{bmatrix}T^{-1}\right) = T\operatorname{sign}\left(\begin{bmatrix}A & 0\\ 0 & -A^{T}\end{bmatrix}\right)T^{-1}$$
$$= \begin{bmatrix}-I_{n} & 2X\\ 0 & I_{n}\end{bmatrix}.$$

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Apply sign function iteration
$$Z \leftarrow \frac{1}{2}(Z + Z^{-1})$$
 to $H = \begin{bmatrix} A & W \\ 0 & -A^T \end{bmatrix}$:

$$H + H^{-1} = \begin{bmatrix} A & W \\ 0 & -A^T \end{bmatrix} + \begin{bmatrix} A^{-1} & A^{-1}WA^{-T} \\ 0 & -A^{-T} \end{bmatrix}$$

 \implies Sign function iteration for Lyapunov equation:

$$\begin{array}{ll} \mathcal{A}_{0} \leftarrow \mathcal{A}, & \mathcal{A}_{j+1} \leftarrow \frac{1}{2} \left(\mathcal{A}_{j} + \mathcal{A}_{j}^{-1} \right), \\ \mathcal{W}_{0} \leftarrow \mathcal{G}, & \mathcal{W}_{j+1} \leftarrow \frac{1}{2} \left(\mathcal{W}_{j} + \mathcal{A}_{j}^{-1} \mathcal{W}_{j} \mathcal{A}_{j}^{-T} \right), \end{array} \qquad j = 0, 1, 2, \dots.$$

Define $A_{\infty} := \lim_{j \to \infty} A_j$, $W_{\infty} := \lim_{j \to \infty} W_j$.

Theorem

If A is Hurwitz, then

$$A_{\infty} = -I_n$$
 and $X = \frac{1}{2}W_{\infty}$.

Factored form

Recall sign function iteration for $AX + XA^T + W = 0$:

$$\begin{array}{ll} \mathcal{A}_0 \leftarrow \mathcal{A}, & \mathcal{A}_{j+1} \leftarrow \frac{1}{2} \left(\mathcal{A}_j + \mathcal{A}_j^{-1} \right), \\ \mathcal{W}_0 \leftarrow \mathcal{G}, & \mathcal{W}_{j+1} \leftarrow \frac{1}{2} \left(\mathcal{W}_j + \mathcal{A}_j^{-1} \mathcal{W}_j \mathcal{A}_j^{-T} \right), \end{array} \qquad j = 0, 1, 2, \dots.$$

Now consider the second iteration for $W = BB^{T}$, starting with $W_0 = BB^{T} =: B_0 B_0^{T}:$

$$\frac{1}{2} \left(W_j + A_j^{-1} W_j A_j^{-T} \right) = \frac{1}{2} \left(B_j B_j^T + A_j^{-1} B_j B_j^T A_j^{-T} \right)$$

$$= \frac{1}{2} \left[B_j \quad A_j^{-1} B_j \right] \left[B_j \quad A_j^{-1} B_j \right]^T$$

Hence, obtain factored iteration

$$B_{j+1} \leftarrow \frac{1}{\sqrt{2}} \begin{bmatrix} B_j & A_j^{-1}B_j \end{bmatrix}$$

with
$$S := \frac{1}{\sqrt{2}} \lim_{j \to \infty} B_j$$
 and $X = SS^T$.

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$$\frac{1}{2} \left(W_j + A_j^{-1} W_j A_j^{-T} \right) = \frac{1}{2} \left(B_j B_j^T + A_j^{-1} B_j B_j^T A_j^{-T} \right)$$

= $\frac{1}{2} \left[B_j \quad A_j^{-1} B_j \right] \left[B_j \quad A_j^{-1} B_j \right]^T .$

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Factored form

Factored sign function iteration for $A(SS^{T}) + (SS^{T})A^{T} + BB^{T} = 0$

$$\begin{array}{ll} A_0 \leftarrow A, & A_{j+1} \leftarrow \frac{1}{2} \left(A_j + A_j^{-1} \right), \\ B_0 \leftarrow B, & B_{j+1} \leftarrow \frac{1}{\sqrt{2}} \left[B_j & A_j^{-1} B_j \right], \end{array} \qquad j = 0, 1, 2, \dots .$$

Remarks:

• To get both Gramians, run in parallel

$$\mathcal{C}_{j+1} \ \leftarrow \ rac{1}{\sqrt{2}} \left[egin{matrix} \mathcal{C}_{j} \ \mathcal{C}_{j} \mathcal{A}_{j}^{-1} \end{bmatrix}.$$

- To avoid growth in numbers of columns of B_j (or rows of C_j): column compression by RRLQ or truncated SVD.
- Several options to incorporate scaling, e.g., scale "A"-iteration only.
- Simple stopping criterion: $||A_j + I_n||_F \leq tol$.

Numerical Methods for Solving Lyapunov Equations The ADI Method

Recall Peaceman Rachford ADI:

Consider Au = s where $A \in \mathbb{R}^{n \times n}$ spd, $s \in \mathbb{R}^n$. ADI Iteration Idea: Decompose A = H + V with $H, V \in \mathbb{R}^{n \times n}$ such that

$$(H + pI)v = r$$
$$(V + pI)w = t$$

can be solved easily/efficiently.

ADI Iteration

If $H, V \text{ spd} \Rightarrow \exists p_k, k = 1, 2, \dots$ such that

$$u_{0} = 0$$

(H + p_{k}l)u_{k-\frac{1}{2}} = (p_{k}l - V)u_{k-1} + s
(V + p_{k}l)u_{k} = (p_{k}l - H)u_{k-\frac{1}{2}} + s

converges to $u \in \mathbb{R}^n$ solving Au = s.

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Numerical Methods for Solving Lyapunov Equations

The Lyapunov operator

$$\mathcal{L}: X \mapsto AX + XA^T$$

can be decomposed into the linear operators

 $\mathcal{L}_H: X \mapsto AX, \qquad \mathcal{L}_V: X \mapsto XA^T.$

In analogy to the standard ADI method we find the



Numerical Methods for Solving Lyapunov Equations Low-Rank ADI

Consider $AX + XA^T = -BB^T$ for stable A; $B \in \mathbb{R}^{n \times m}$ with $m \ll n$.

[Wachspress '95]

For $k = 1, \ldots, k_{\max}$

$$\begin{array}{rcl} X_0 & = & 0 \\ (A+p_kI)X_{k-\frac{1}{2}} & = & -BB^T - X_{k-1}(A^T - p_kI) \\ (A+p_kI)X_k^T & = & -BB^T - X_{k-\frac{1}{2}}^T(A^T - p_kI) \end{array}$$

Rewrite as one step iteration and factorize $X_k = Z_k Z_k^T$, $k = 0, \ldots, k_{max}$

$$Z_{0}Z_{0}^{T} = 0$$

$$Z_{k}Z_{k}^{T} = -2p_{k}(A + p_{k}I)^{-1}BB^{T}(A + p_{k}I)^{-T} + (A + p_{k}I)^{-1}(A - p_{k}I)Z_{k-1}Z_{k-1}^{T}(A - p_{k}I)^{T}(A + p_{k}I)^{-T}$$

 $\ldots \rightsquigarrow$ low-rank Cholesky factor ADI

[PENZL '97/'00, LI/WHITE '99/'02, B./LI/PENZL '99/'08, GUGERCIN/SORENSEN/ANTOULAS '03]

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... ~> low-rank Cholesky factor ADI

[PENZL '97/'00, LI/WHITE '99/'02, B./LI/PENZL '99/'08, GUGERCIN/SORENSEN/ANTOULAS '03]
Solving Large-Scale Matrix Equations Numerical Methods for Solving Lyapunov Equations

$$Z_k = [\sqrt{-2p_k}(A + p_kI)^{-1}B, (A + p_kI)^{-1}(A - p_kI)Z_{k-1}]$$

[PENZL '00]

Observing that $(A - p_i I)$, $(A + p_k I)^{-1}$ commute, we rewrite $Z_{k_{\max}}$ as

$$Z_{k_{\max}} = [z_{k_{\max}}, P_{k_{\max}-1}z_{k_{\max}}, P_{k_{\max}-2}(P_{k_{\max}-1}z_{k_{\max}}), \dots, P_1(P_2 \cdots P_{k_{\max}-1}z_{k_{\max}})],$$

[LI/WHITE '02]

where

$$z_{k_{\max}} = \sqrt{-2p_{k_{\max}}}(A + p_{k_{\max}}I)^{-1}B$$

and

$$P_i := \frac{\sqrt{-2p_i}}{\sqrt{-2p_{i+1}}} \left[I - (p_i + p_{i+1})(A + p_i I)^{-1} \right].$$

Solving Large-Scale Matrix Equations Numerical Methods for Solving Lyapunov Equations

$$Z_{k} = \left[\sqrt{-2p_{k}}(A + p_{k}I)^{-1}B, \ (A + p_{k}I)^{-1}(A - p_{k}I)Z_{k-1}\right]$$

[PENZL '00]

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Lyapunov equation $0 = AX + XA^T + BB^T$.

Algorithm [Penzl '97/'00, LI/WHITE '99/'02, B. 04, B./LI/PENZL '99/'08]

$$V_{1} \leftarrow \sqrt{-2 \operatorname{re} p_{1}} (A + p_{1}I)^{-1}B, \quad Z_{1} \leftarrow V_{1}$$

FOR $k = 2, 3, ...$

$$V_{k} \leftarrow \sqrt{\frac{\operatorname{re} p_{k}}{\operatorname{re} p_{k-1}}} (V_{k-1} - (p_{k} + \overline{p_{k-1}})(A + p_{k}I)^{-1}V_{k-1})$$

$$Z_{k} \leftarrow [Z_{k-1} \quad V_{k}]$$

$$Z_{k} \leftarrow \operatorname{rrlq}(Z_{k}, \tau) \quad \text{column compression}$$

At convergence, $Z_{k_{\max}} Z_{k_{\max}}^T \approx X$, where (without column compression)

$$Z_{k_{\max}} = \begin{bmatrix} V_1 & \dots & V_{k_{\max}} \end{bmatrix}, \quad V_k = \begin{bmatrix} \mathbb{C}^{n \times m} \end{bmatrix}$$

Note: Implementation in real arithmetic possible by combining two steps [B./Li/Penzl '99/'08] or using new idea employing the relation of 2 consecutive complex factors [B./Kürschner/Saak '11].

Numerical Methods for Solving Lyapunov Equations Lyapunov equation $0 = AX + XA^T + BB^T$.

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Numerical Results for ADI Optimal Cooling of Steel Profiles

• Mathematical model: boundary control for linearized 2D heat equation.

$$c \cdot \rho \frac{\partial}{\partial t} x = \lambda \Delta x, \quad \xi \in \Omega$$

$$\lambda \frac{\partial}{\partial n} x = \kappa (u_k - x), \quad \xi \in \Gamma_k, \ 1 \le k \le 7$$

$$\frac{\partial}{\partial n} x = 0, \qquad \xi \in \Gamma_7.$$

$$\implies m = 7, q = 6.$$

- FEM Discretization, different models for initial mesh (n = 371),
 1, 2, 3, 4 steps of mesh refinement ⇒
 - n = 1357, 5177, 20209, 79841.



Source: Physical model: courtesy of Mannesmann/Demag.

Math. model: Tröltzsch/Unger 1999/2001, Penzl 1999, Saak 2003.

Numerical Results for ADI Optimal Cooling of Steel Profiles

• Solve dual Lyapunov equations needed for balanced truncation, i.e.,

 $APM^{T} + MPA^{T} + BB^{T} = 0, \quad A^{T}QM + M^{T}QA + C^{T}C = 0,$

for n = 79,841.

- 25 shifts chosen by Penzl heuristic from 50/25 Ritz values of A of largest/smallest magnitude, no column compression performed.
- No factorization of mass matrix required.
- Computations done on Core2Duo at 2.8GHz with 3GB RAM and 32Bit-MATLAB.



Numerical Results for ADI

Computations by Martin Köhler '10

- A ∈ ℝ^{n×n} ≡ FDM matrix for 2D heat equation on [0, 1]² (LYAPACK benchmark demo_l1, m = 1).
- 16 shifts chosen by Penzl heuristic from 50/25 Ritz values of A of largest/smallest magnitude.
- Computations on 2 dual core Intel Xeon 5160 with 16 GB RAM using M.E.S.S. (http://svncsc.mpi-magdeburg.mpg.de/trac/messtrac/).

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n	M.E.S.S. ¹ (C)	LyaPack	M.E.S.S. (MATLAB)
100	0.023	0.124	0.158
625	0.042	0.104	0.227
2,500	0.159	0.702	0.989
10,000	0.965	6.22	5.644
40,000	11.09	71.48	34.55
90,000	34.67	418.5	90.49
160,000	109.3	out of memory	219.9
250,000	193.7	out of memory	403.8
562,500	930.1	out of memory	1216.7
1,000,000	2220.0	out of memory	2428.6

CPU Times

Numerical Results for ADI

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Projection-based methods for Lyapunov equations with $A + A^T < 0$:

- Compute orthonormal basis range (Z), Z ∈ ℝ^{n×r}, for subspace Z ⊂ ℝⁿ, dim Z = r.
- **③** Solve small-size Lyapunov equation $\hat{A}\hat{X} + \hat{X}\hat{A}^{T} + \hat{B}\hat{B}^{T} = 0$.
- Use $X \approx Z \hat{X} Z^T$.

Examples:

• Krylov subspace methods, i.e., for m = 1:

$$\mathcal{Z} = \mathcal{K}(A, B, r) = \operatorname{span}\{B, AB, A^2B, \dots, A^{r-1}B\}$$

[SAAD '90, JAIMOUKHA/KASENALLY '94, JBILOU '02-'08].

• K-PIK [Simoncini '07],

$$\mathcal{Z} = \mathcal{K}(A, B, r) \cup \mathcal{K}(A^{-1}, B, r).$$

• Rational Krylov [DRUSKIN/SIMONCINI '11].

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Examples:

• ADI subspace [B./R.-C. LI/TRUHAR '08]:

$$\mathcal{Z} = \operatorname{colspan} \left[\begin{array}{cc} V_1, & \dots, & V_r \end{array} \right].$$

Note:

- ADI subspace is rational Krylov subspace [J.-R. LI/WHITE '02].
- Similar approach: ADI-preconditioned global Arnoldi method [JBILOU '08].

Numerical Methods for Solving Lyapunov Equations Numerical examples for Galerkin-ADI

FEM semi-discretized control problem for parabolic PDE:

- optimal cooling of rail profiles,
- n = 20, 209, m = 7, q = 6.



CPU times: 80s (projection every 5th ADI step) vs. 94s (no projection).

Computations by Jens Saak '10.

Numerical Methods for Solving Lyapunov Equations Numerical examples for Galerkin-ADI

FEM semi-discretized control problem for parabolic PDE:

- optimal cooling of rail profiles,
- n = 20, 209, m = 7, q = 6.



CPU times: 368s (projection every 5th ADI step) vs. 1207s (no projection).

Computations by Jens Saak '10.

Numerical Methods for Solving Lyapunov Equations Numerical examples for Galerkin-ADI: optimal cooling of rail profiles, n = 79,841.

M.E.S.S. w/o Galerkin projection and column compression



M.E.S.S. with Galerkin projection and column compression



Solving Large-Scale Matrix Equations Numerical example for BT: Optimal Cooling of Steel Profiles



- BT model computed with sign function method,
- MT w/o static condensation, same order as BT model.

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Solving Large-Scale Matrix Equations Numerical example for BT: Optimal Cooling of Steel Profiles



- BT model computed with sign function method,
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- BT model computed using M.E.S.S. in MATLAB,
- dualcore, computation time:
 <10 min.

Modal Truncation Balanced

Solving Large-Scale Matrix Equations Numerical example for BT: Microgyroscope (Butterfly Gyro)



- By applying AC voltage to electrodes, wings are forced to vibrate in anti-phase in wafer plane.
- Coriolis forces induce motion of wings out of wafer plane yielding sensor data.

- Vibrating micro-mechanical gyroscope for inertial navigation.
- Rotational position sensor.



Source: The Oberwolfach Benchmark Collection http://www.intek.de/simulation/benchmark Courtesy of D. Billger (Imego Institute, Göteborg), Saab Bofors Dynamics AB.

Solving Large-Scale Matrix Equations Numerical example for BT: Microgyroscope (Butterfly Gyro)

• FEM discretization of structure dynamical model using quadratic tetrahedral elements (ANSYS-SOLID187)

 \rightsquigarrow n = 34,722, m = 1, q = 12.

• Reduced model computed using SPARED, r = 30.

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Solving Large-Scale Algebraic Riccati Equations Theory [Lancaster/Rodman '95]

Theorem

Consider the (continuous-time) algebraic Riccati equation (ARE)

$$0 = \mathcal{R}(X) = C^{\mathsf{T}}C + A^{\mathsf{T}}X + XA - XBB^{\mathsf{T}}X,$$

with $A \in \mathbb{R}^{n \times n}$, $B \in \mathbb{R}^{n \times m}$, $C \in \mathbb{R}^{q \times n}$, (A, B) stabilizable, (A, C) detectable. Then:

(a) There exists a unique stabilizing $X_* \in \{X \in \mathbb{R}^{n \times n} | \mathcal{R}(X) = 0\}$, i.e., $\Lambda(A - BB^T X_*) \in \mathbb{C}^-$.

(b)
$$X_* = X_*^T \ge 0$$
 and $X_* \ge X$ for all $X \in \{X \in \mathbb{R}^{n \times n} \mid \mathcal{R}(X) = 0\}$

(c) If (A, C) observable, then $X_* > 0$.

(d) span $\left\{ \begin{bmatrix} I_n \\ -X_* \end{bmatrix} \right\}$ is the unique invariant subspace of the Hamiltonian matrix

$$H = \begin{bmatrix} A & BB^T \\ C^T C & -A^T \end{bmatrix}$$

corresponding to $\Lambda(H) \cap \mathbb{C}^-$.

Solving Large-Scale Algebraic Riccati Equations Numerical Methods [Bini,

[Bini/lannazzo/Meini '12]

Numerical Methods (incomplete list)

- Invariant subspace methods (~> eigenproblem for Hamiltonian matrix):
 - Schur vector method (care)
 - Hamiltonian SR algorithm
 - Symplectic URV-based method

[Laub '79] [Bunse-Gerstner/Mehrmann '86]

[B./Mehrmann/Xu '97/'98, Chu/Liu/Mehrmann '07]

- Spectral projection methods
 - Sign function method [ROBERTS '71, BYERS '87]
 - Disk function method [BAI/DEMMEL/GU '94, B. '97]

(rational, global) Krylov subspace techniques [JAIMOUKHA/KASENALLY '94, JBILOU '03/'06, HEYOUNI/JBILOU '09]

• Newton's method

- Kleinman iteration
- Line search acceleration
- Newton-ADI
- Inexact Newton

[Kleinman '68] [B./Byers '98] [B./J.-R. Li/Penzl '99/'08]

[Feitzinger/Hylla/Sachs '09,B./Heinkenschloss/Saak/Weichelt '15]



Solving Large-Scale Matrix Equations Software

Lyapack

[Penzl 2000]

MATLAB toolbox for solving

- Lyapunov equations and algebraic Riccati equations,
- model reduction and LQR problems.

Main work horse: Low-rank ADI and Newton-ADI iterations.

M.E.S.S. – Matrix Equations Sparse Solvers

[B./Köhler/Saak '08–]

• Extended and revised version of LYAPACK.

 Includes solvers for large-scale differential Riccati equations (based on Rosenbrock and BDF methods).

• Many algorithmic improvements:

- new ADI parameter selection,
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- more efficient use of direct solvers,
- treatment of generalized systems without factorization of the mass matrix,
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Topics Not Covered

- Special methods for second-order (mechanical) systems.
- Extensions to bilinear and stochastic systems.
- Balanced truncation for discrete-time systems.
- Extensions to descriptor systems $E\dot{x} = Ax + Bu$, E singular.
- Frequency-limited/-weighted balanced truncation.
- Application to parametric model reduction:

$$\dot{x} = A(p)x + B(p)u, \quad y = C(p)x,$$

where $p \in \mathbb{R}^d$ is a free parameter vector; parameters should be preserved in the reduced-order model.

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